Random Access Zipper



Presented by Kyle Headley

TFP'16 College Park

Functional programmers want simple data types

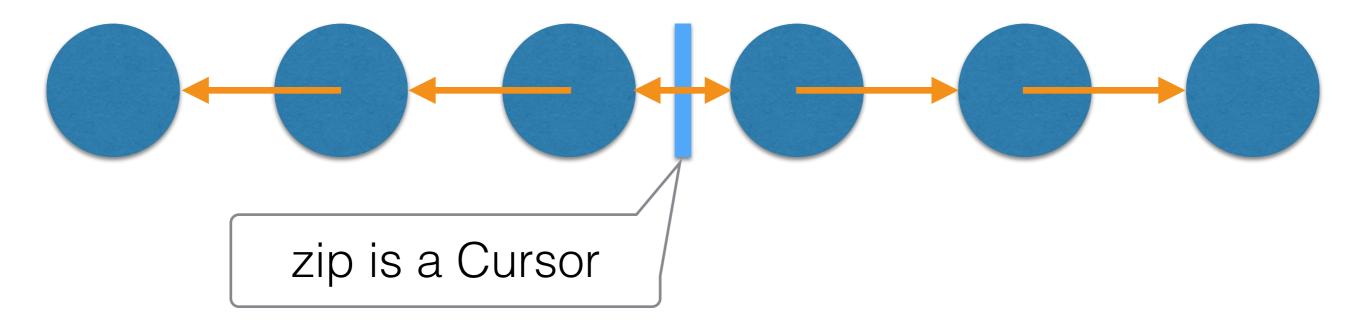
What do we have for sequences?

type 'a list =
| Nil
| Cons of 'a * 'a list
type 'a zip =

'a list * 'a list

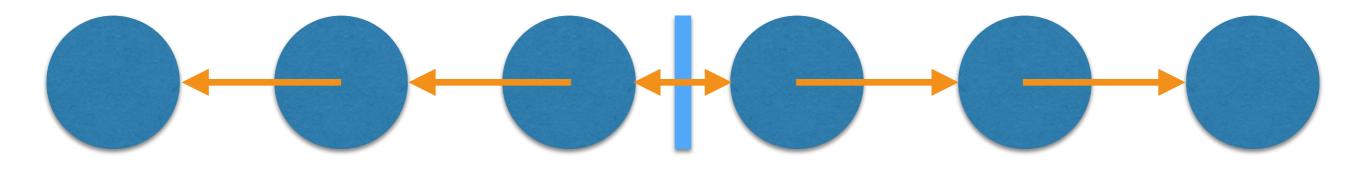
type 'a list = | Nil | Cons of 'a * 'a list

type 'a zip =
 'a list * 'a list



type 'a list =
| Nil
| Cons of 'a * 'a list

type 'a zip =
 'a list * 'a list



Zippers are great

type 'a list =
| Nil
| Cons of 'a * 'a list

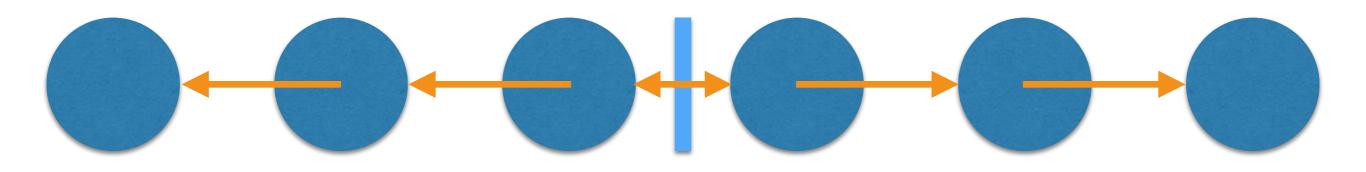
type 'a zip =
 'a list * 'a list

All O(1)!

Zippers are great

type 'a list =
| Nil
| Cons of 'a * 'a list

type 'a zip =
 'a list * 'a list

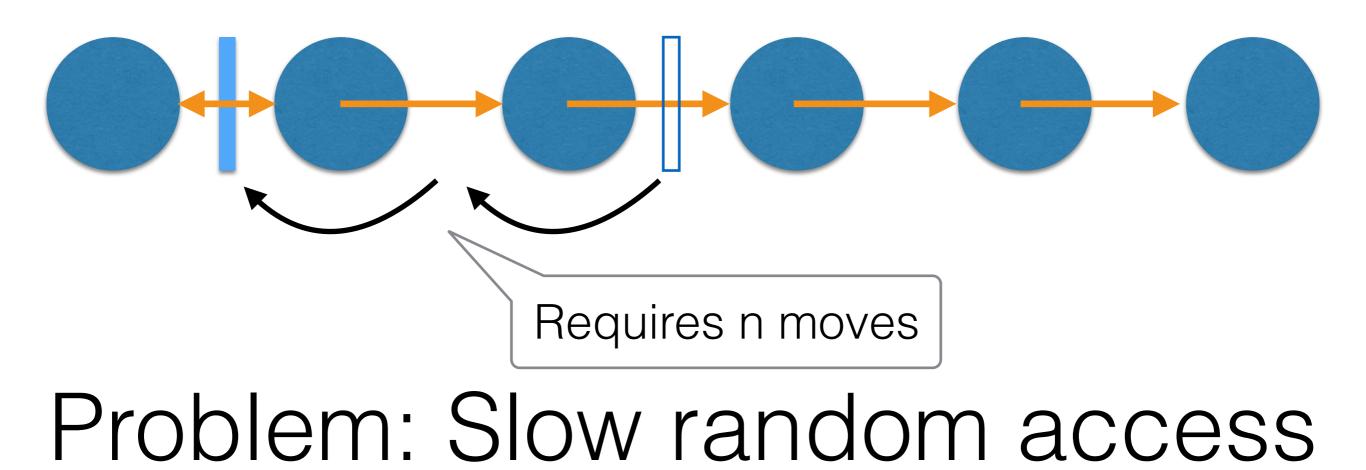


Problem: Slow random access

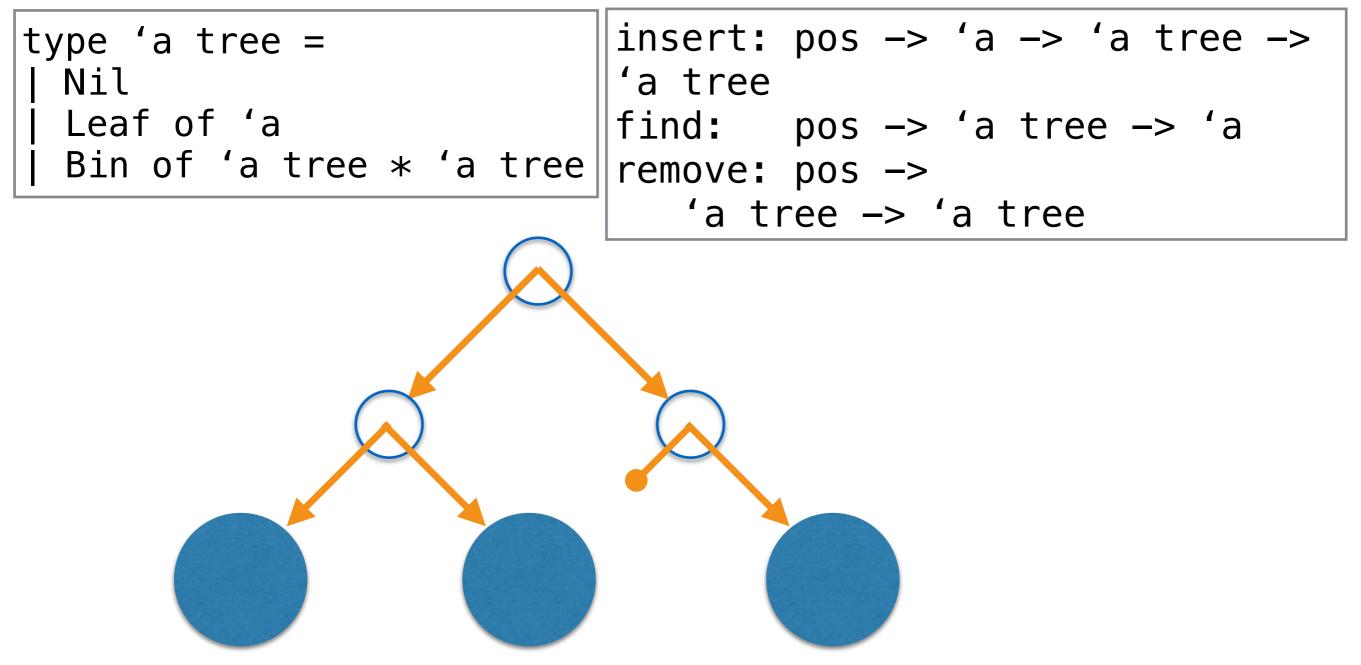
Zippers are great

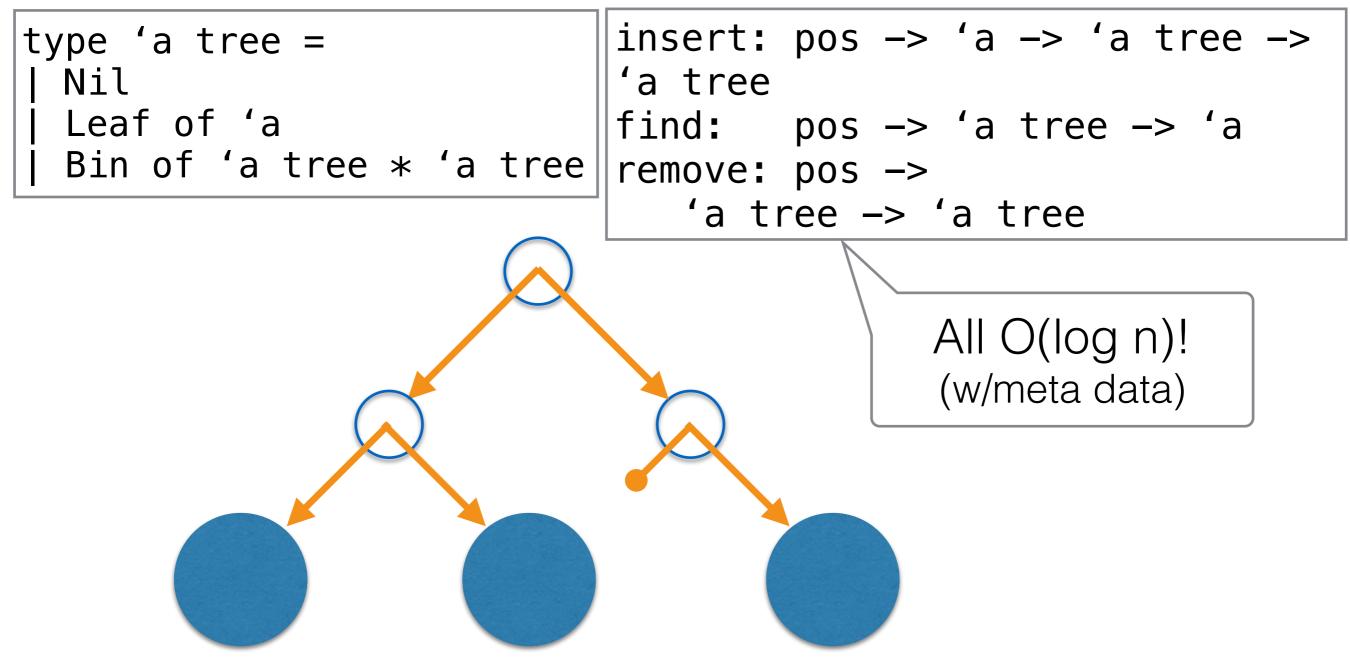
type 'a list =
| Nil
| Cons of 'a * 'a list

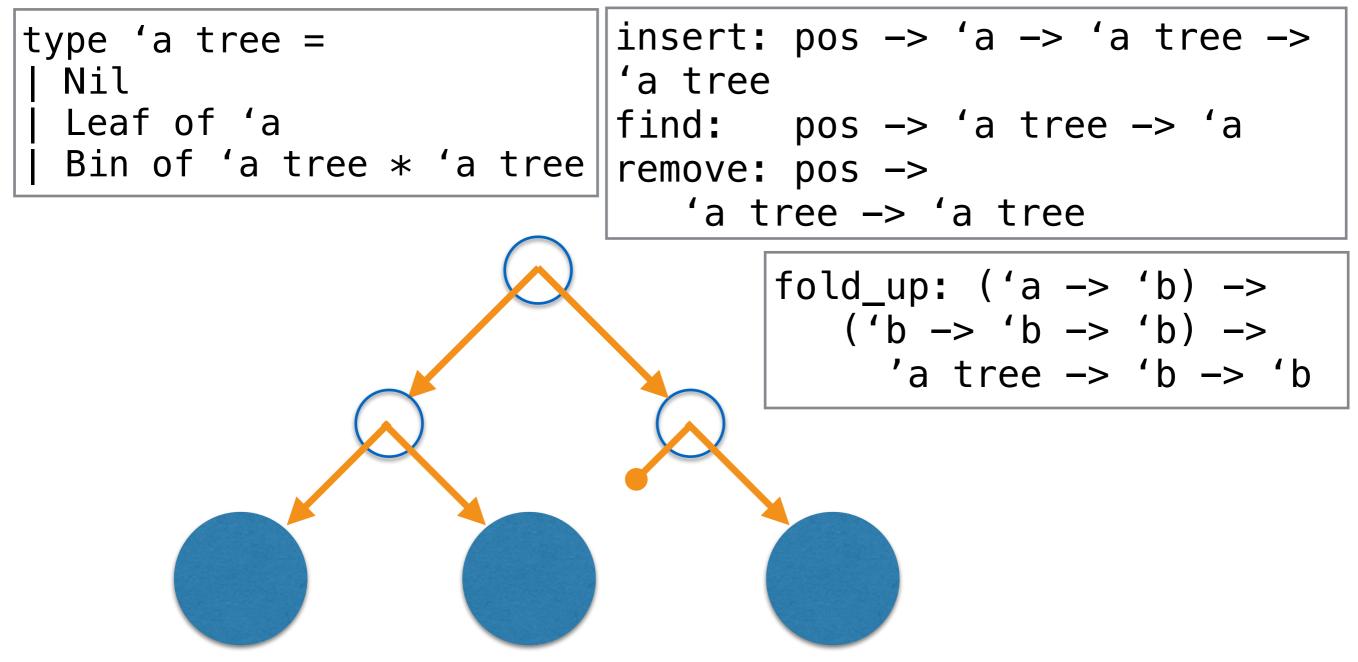
type 'a zip =
 'a list * 'a list

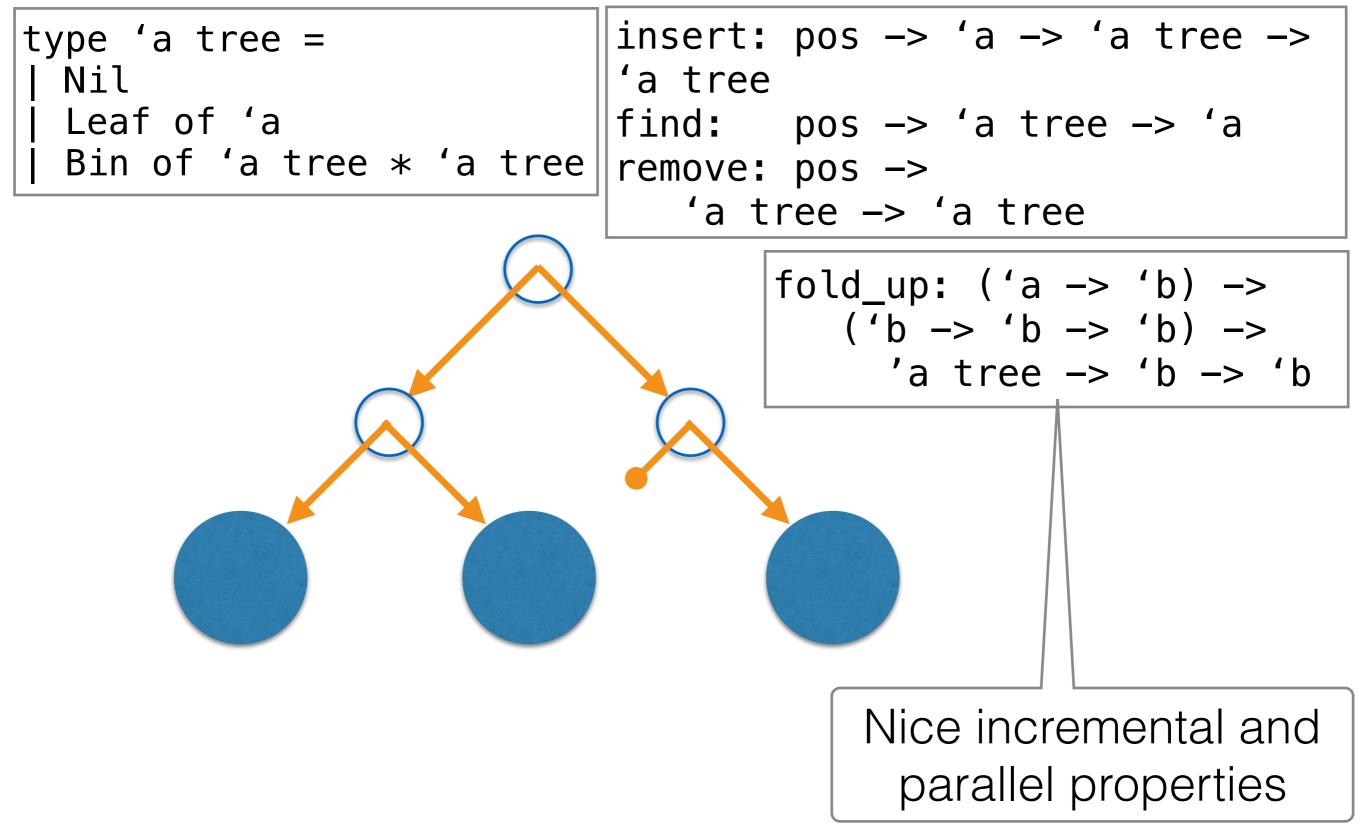


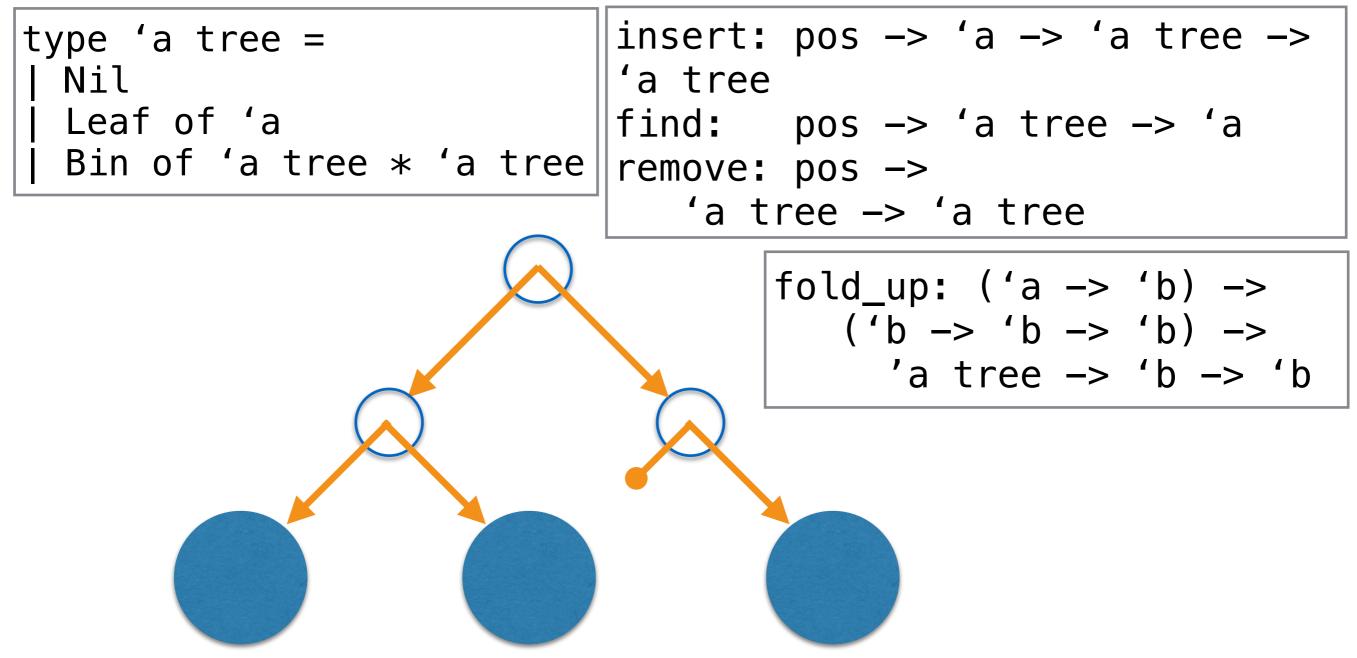
type 'a tree =
| Nil
| Leaf of 'a
| Bin of 'a tree * 'a tree



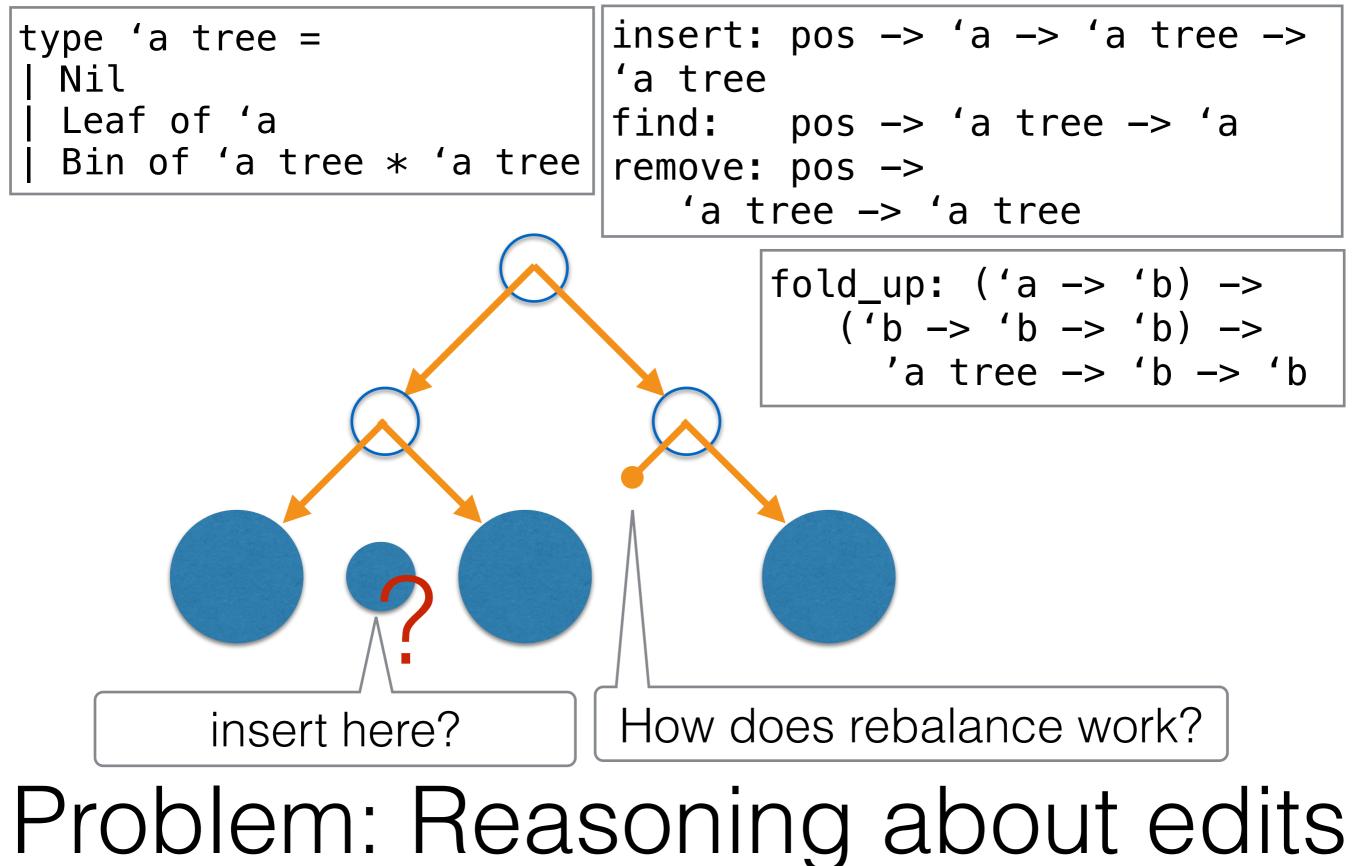






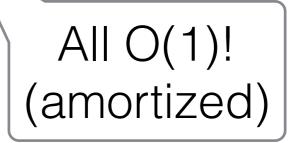


Problem: Reasoning about edits



first: 'a finger -> 'a
last: 'a finger -> 'a
cons: 'a ->
 'a finger -> 'a finger
snoc: 'a ->
 'a finger -> 'a finger

first: 'a finger -> 'a
last: 'a finger -> 'a
cons: 'a ->
 'a finger -> 'a finger
snoc: 'a ->
 'a finger -> 'a finger



first: 'a finger -> 'a
last: 'a finger -> 'a
cons: 'a ->
 'a finger -> 'a finger
snoc: 'a ->
 'a finger -> 'a finger

split: pos -> 'a finger ->
 ('a finger, 'a finger)
append: 'a finger -> 'a finger
-> 'a finger -> 'a finger

first: 'a finger -> 'a
last: 'a finger -> 'a
cons: 'a ->
 'a finger -> 'a finger
snoc: 'a ->
 'a finger -> 'a finger

split: pos -> 'a finger -> ('a finger, 'a finger) append: 'a finger -> 'a finger -> 'a finger -> 'a finger

Both O(log n)!

first: 'a finger -> 'a
last: 'a finger -> 'a
cons: 'a ->
 'a finger -> 'a finger
snoc: 'a ->
 'a finger -> 'a finger

split: pos -> 'a finger ->
 ('a finger, 'a finger)
append: 'a finger -> 'a finger
-> 'a finger -> 'a finger

Problem: Not so simple

Fingertrees are great

```
type 'a node =
 Node2 of 'a * `a
| Node3 of 'a * `a * `a
type 'a digit =
 One of 'a
 Two of 'a * 'a
| Three of 'a * 'a * 'a
| Four of 'a * 'a * 'a * 'a
type 'a finger =
 Nil
 Single of 'a
 Deep of
  'a digit
* ('a node) finger
* 'a digit
```

first:	'a finger −> 'a
last:	'a finger −> 'a
cons:	'a ->
'a	finger -> 'a finger
snoc:	'a ->
'a	finger -> 'a finger

```
split: pos -> 'a finger ->
  ('a finger, 'a finger)
append: 'a finger -> 'a finger
-> 'a finger -> 'a finger
```

Problem: Not so simple

Fingertrees are great

type 'a node = Node2 of 'a * `a | Node3 of 'a \ast `a \ast `a type 'a digit = One of 'a Two of 'a * 'a | Three of 'a * 'a * 'a | Four of 'a * 'a * 'a * 'a type 'a finger = Nil Single of 'a Deep of 'a digit * ('a node) finger * 'a digit

first: 'a finger -> 'a
last: 'a finger -> 'a
cons: 'a ->
 'a finger -> 'a finger
snoc: 'a ->
 'a finger -> 'a finger

split: pos -> 'a finger ->
 ('a finger, 'a finger)
append: 'a finger -> 'a finger
-> 'a finger -> 'a finger

Nested type

Problem: Not so simple

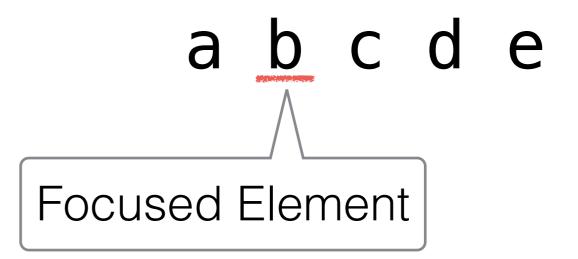
Alternative: Random Access Zipper

Accessible
Editable
Simple

raz

a b c d e

raz



raz |> insert left n a n b c d e

raz |> insert left n a n b c d e |> remove left a b c d e

a b c d e

raz |> insert left n |> remove left |> remove right

a b c d e a n b c d e a b c d e a b d e

raz |> insert left n |> remove left |> remove right |> unfocus

a b c d e a b c d e a b c d e a b d e a b d e

raz |> insert left n |> remove left |> remove right |> unfocus |> focus 0

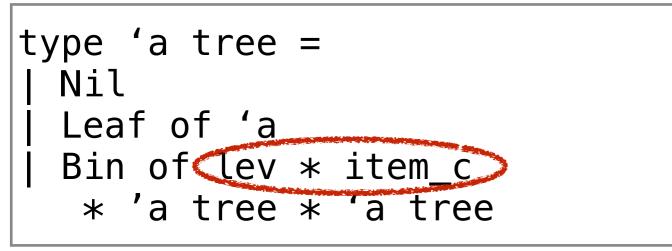
a b c d e a n b c d e a b c d e a b d e a b d e a b d e

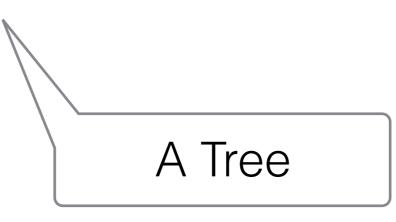
raz > insert left n > remove left > remove right > unfocus > focus 0 > alter right n

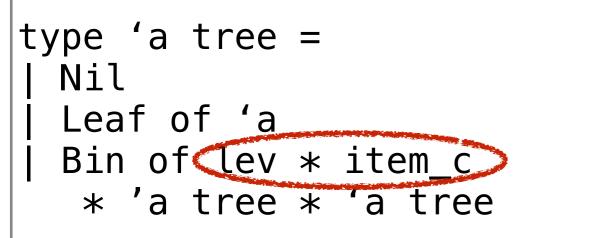
a b c d e a n <u>b</u> c d e a b c d e a b d e a b d e a b d e a n d e

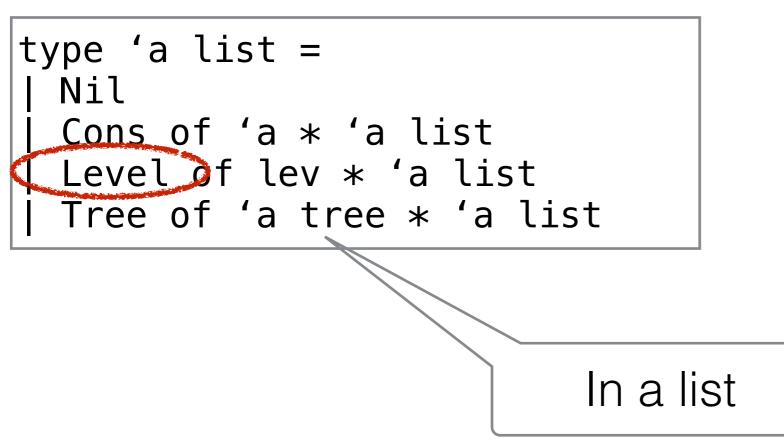
The RAZ is great

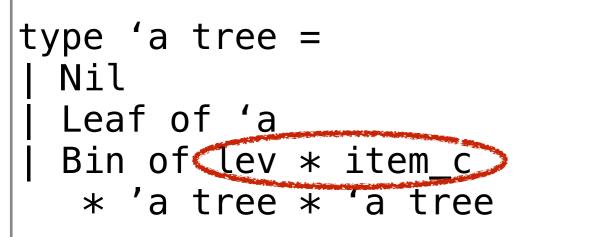
The RAZ is great

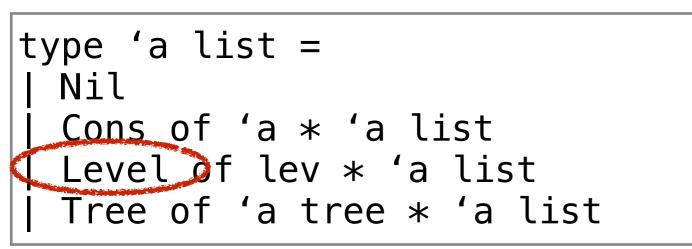












type 'a raz =
 'a list * 'a * 'a list

As a zipper

type 'a tree =
| Nil
| Leaf of 'a
| Bin of lev * item_c
 * 'a tree * 'a tree

type 'a list =
| Nil
| Cons of 'a * 'a list
| Level of lev * 'a list
| Tree of 'a tree * 'a list

type 'a raz =
 'a list * 'a * 'a list

fold_up: ('a -> 'b) ->
 ('b -> 'b -> 'b) ->
 'a tree -> 'b -> 'b

Still get tree info

type 'a tree =
| Nil
| Leaf of 'a
| Bin of lev * item_c
 * 'a tree * 'a tree

type 'a list = | Nil | Cons of 'a * 'a list | Level of lev * 'a list | Tree of 'a tree * 'a list

type 'a raz =
 'a list * 'a * 'a list

fold_up: ('a -> 'b) ->

('b -> 'b -> 'b) ->

'a tree -> 'b -> 'b

All O(1)!

type 'a tree =
| Nil
| Leaf of 'a
| Bin of lev * item_c
 * 'a tree * 'a tree

type 'a list = | Nil | Cons of 'a * 'a list | Level of lev * 'a list | Tree of 'a tree * 'a list

type 'a raz =
 'a list * 'a * 'a list

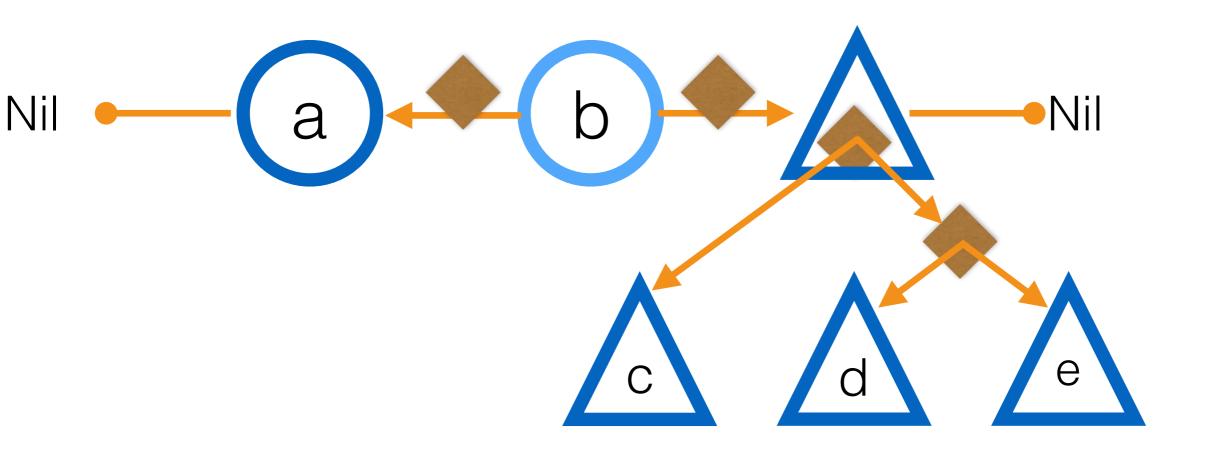
fold_up: ('a -> 'b) ->

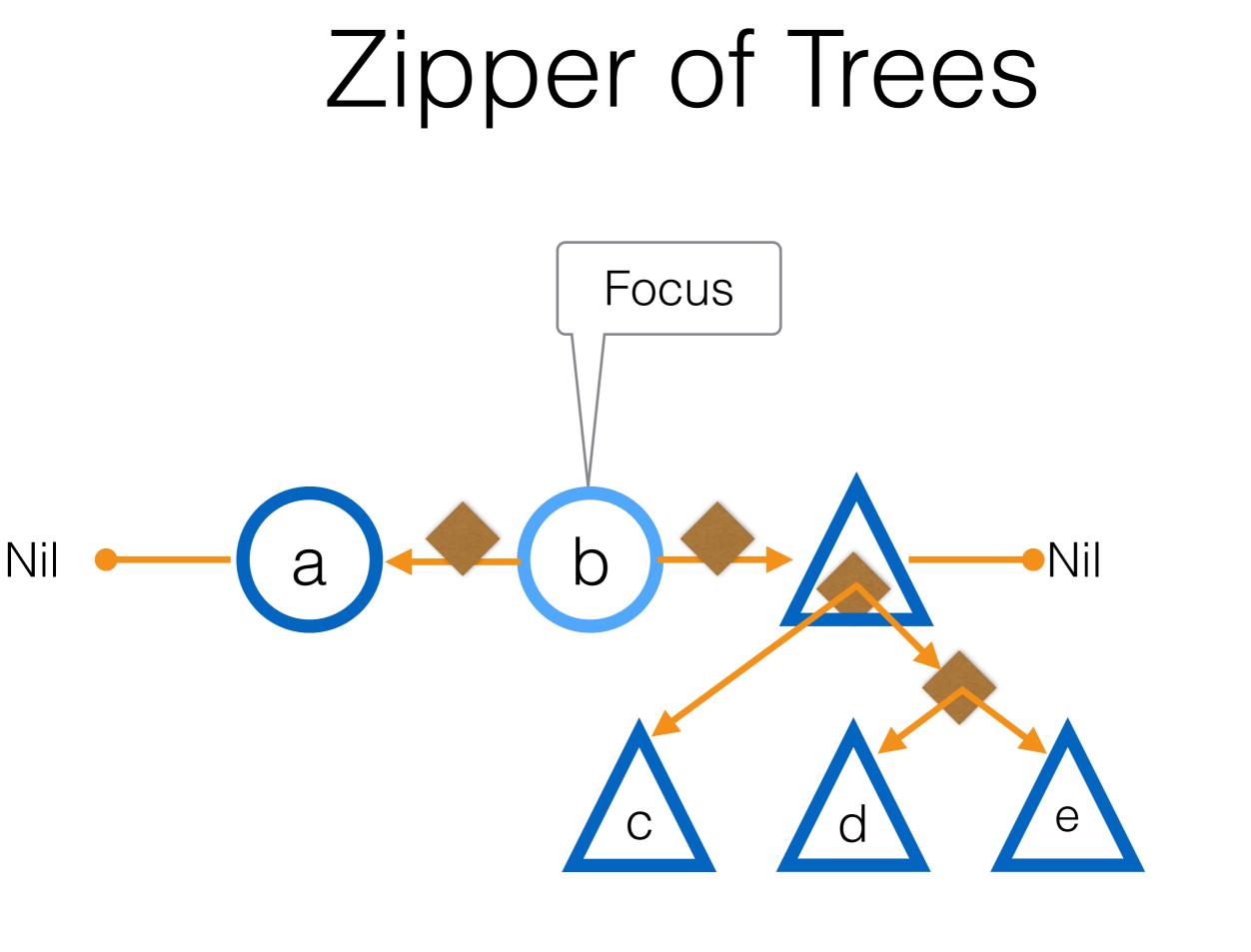
('b -> 'b -> 'b) ->

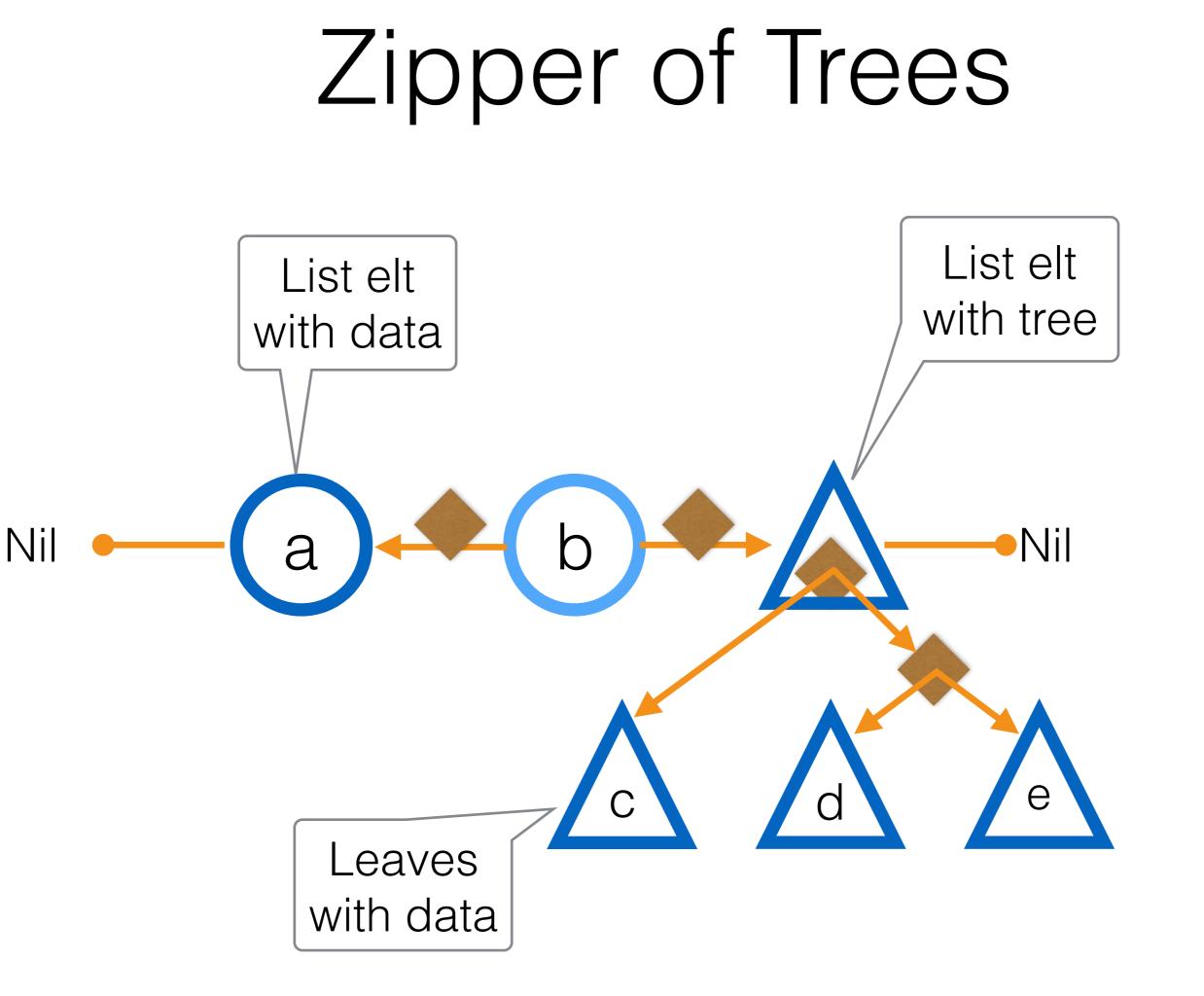
'a tree -> 'b -> 'b

> Both O(log n)! (plus net insertions)

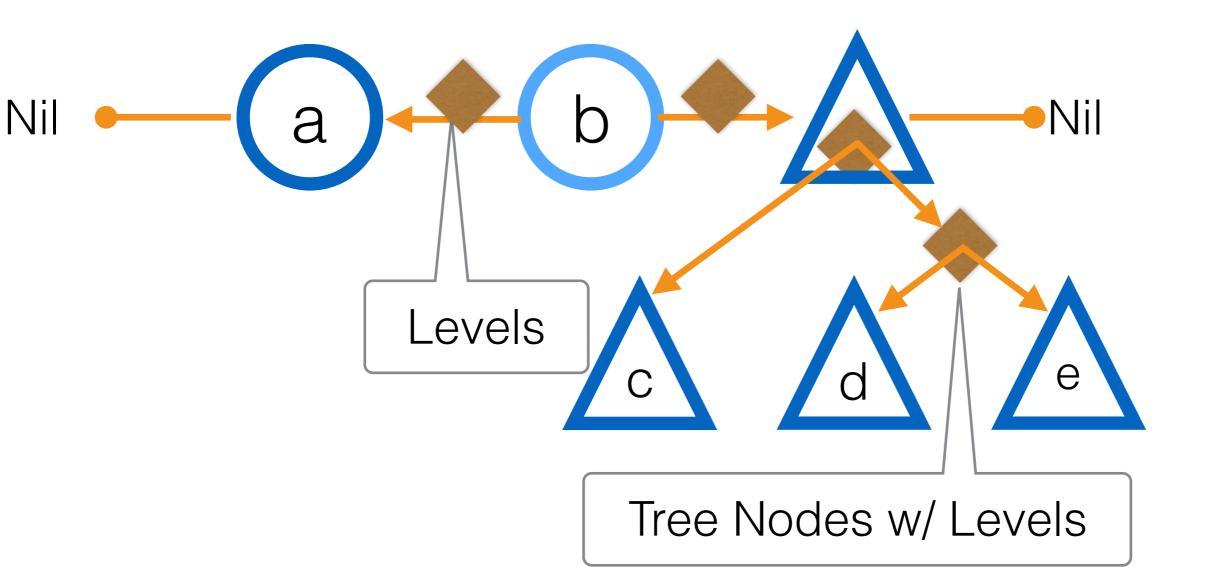
Zipper of Trees







Zipper of Trees



Balance

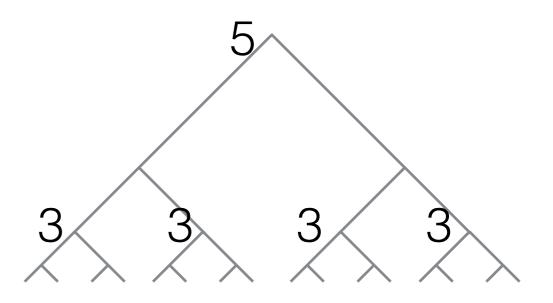
Balance

We use a probabilistic balance, inserting random numbers as levels Because of the way randomness behaves, we get good balance at scale

Balance

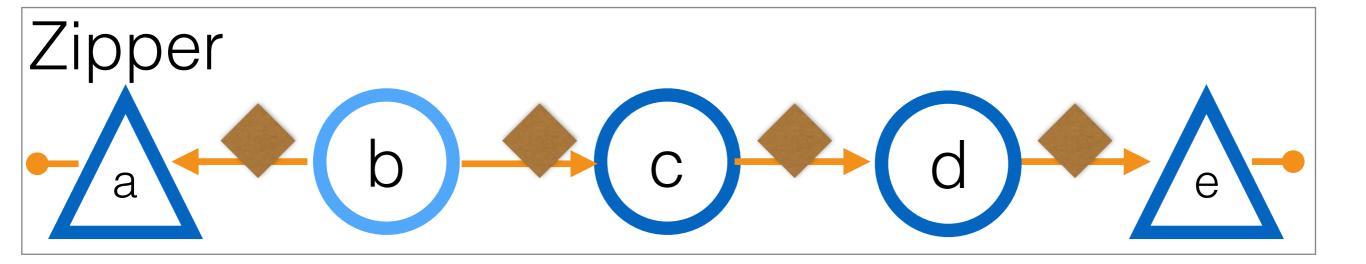
We use a probabilistic balance, inserting random numbers as levels Because of the way randomness behaves, we get good balance at scale

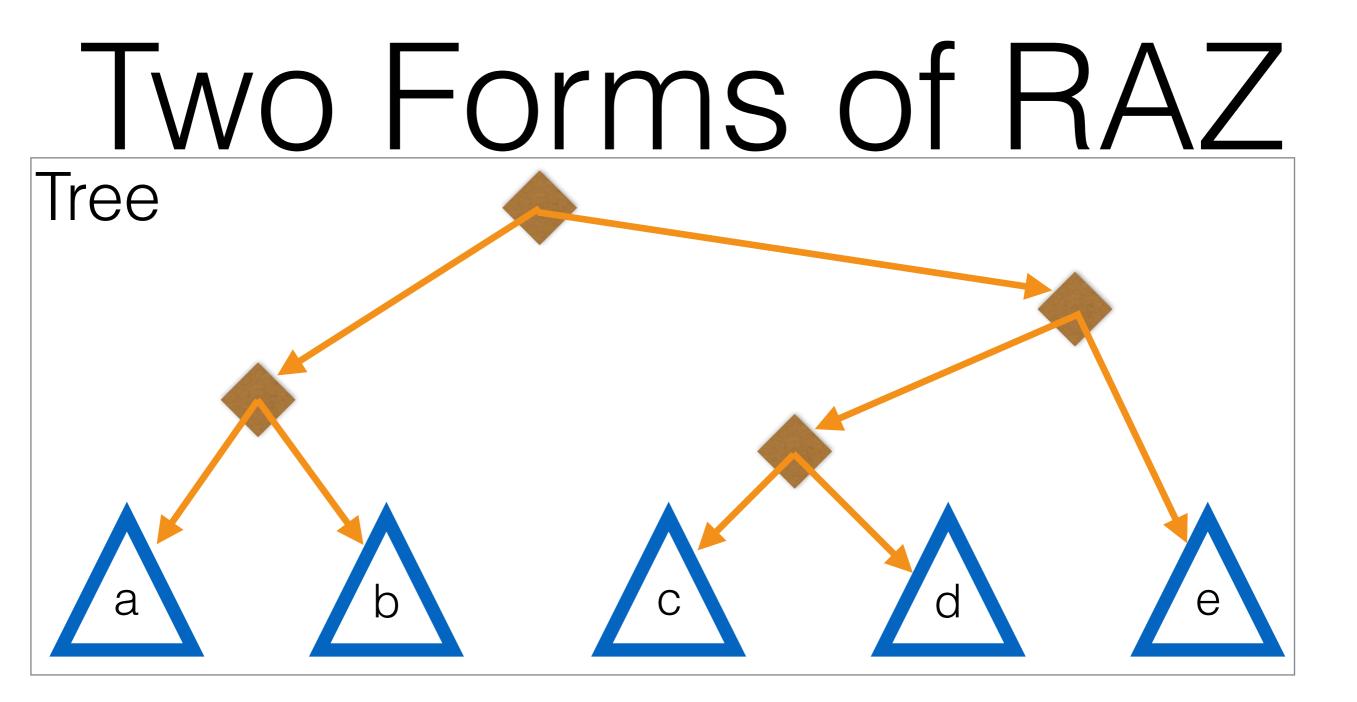
Choose a random level based on balanced tree height distribution

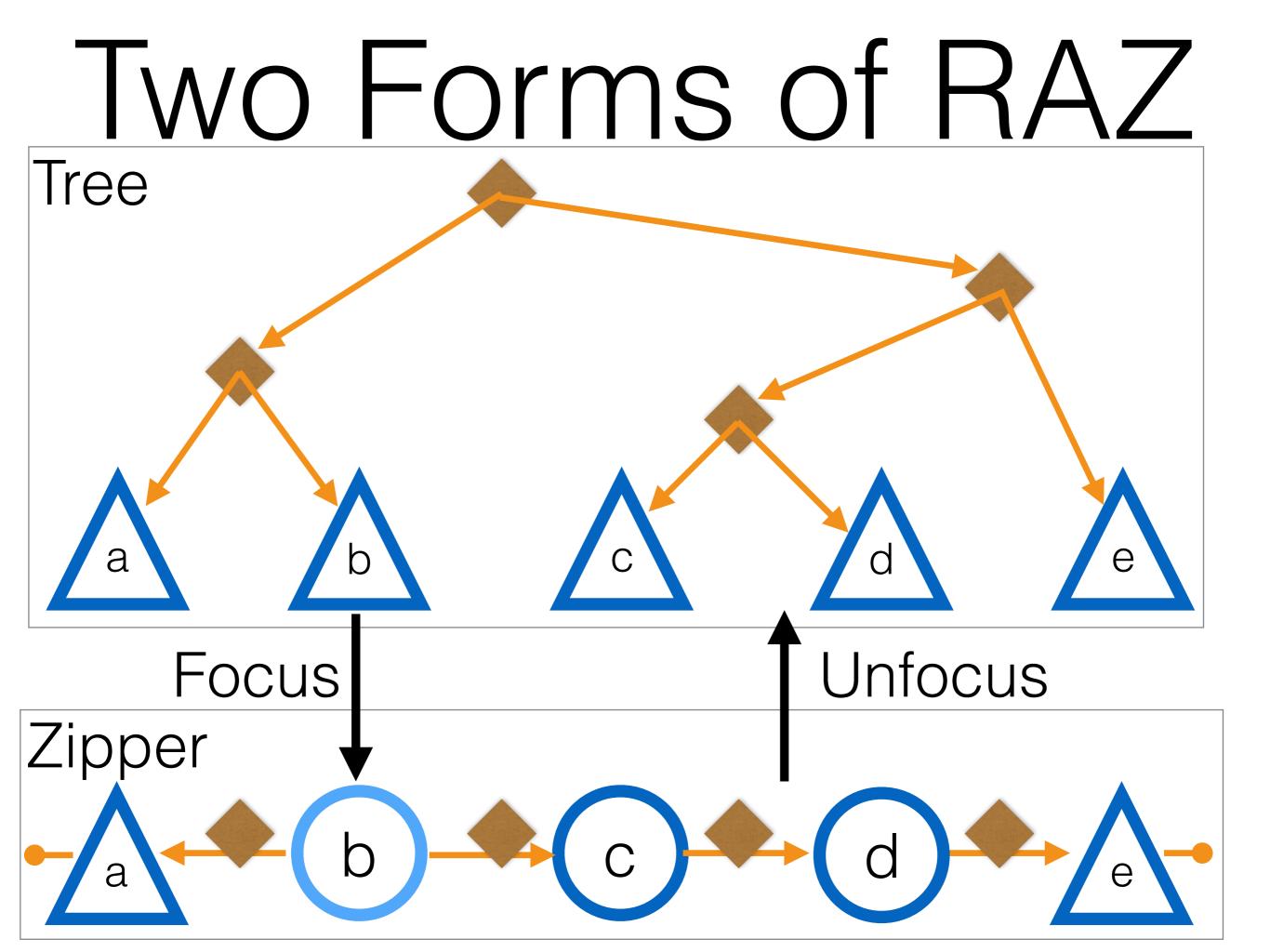


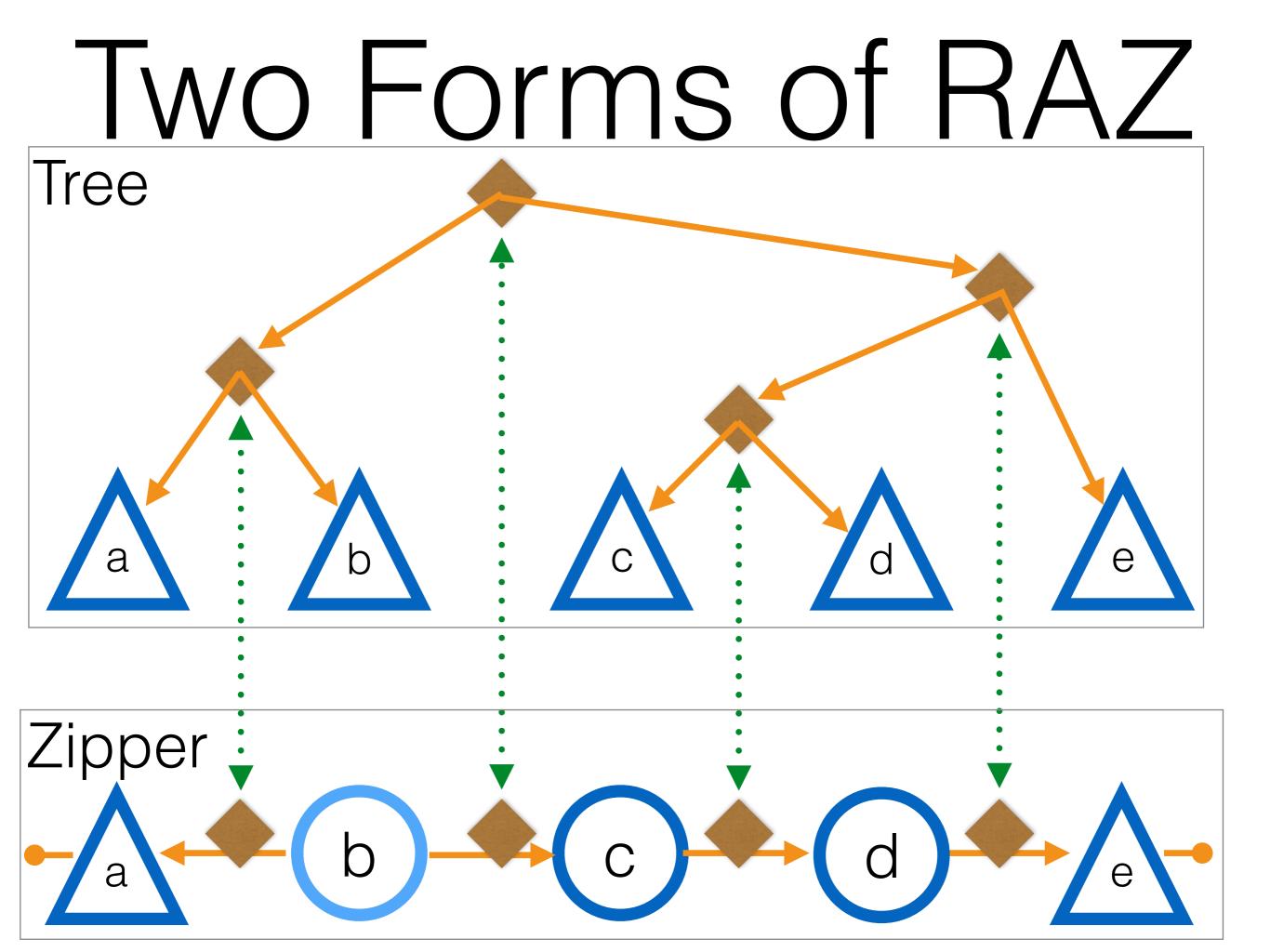
Two Forms of RAZ

Two Forms of RAZ



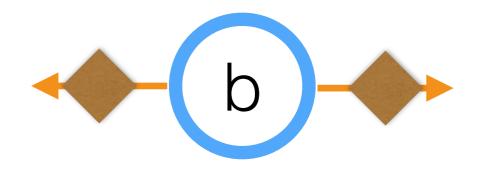






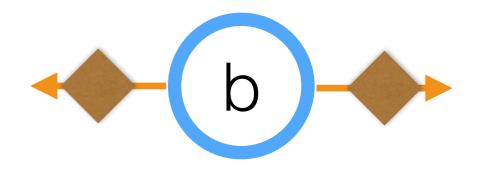
Invariants

Levels on each side of the focused element

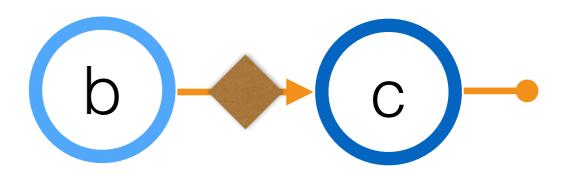


Invariants

Levels on each side of the focused element



Levels between each element except Nil

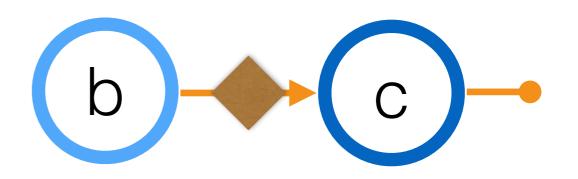


Invariants

Levels on each side of the focused element

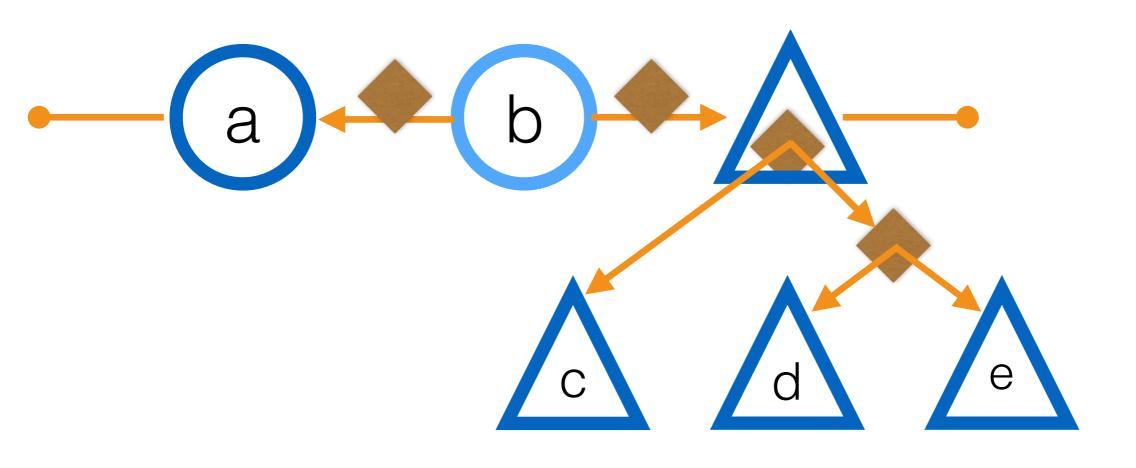


Levels between each element except Nil

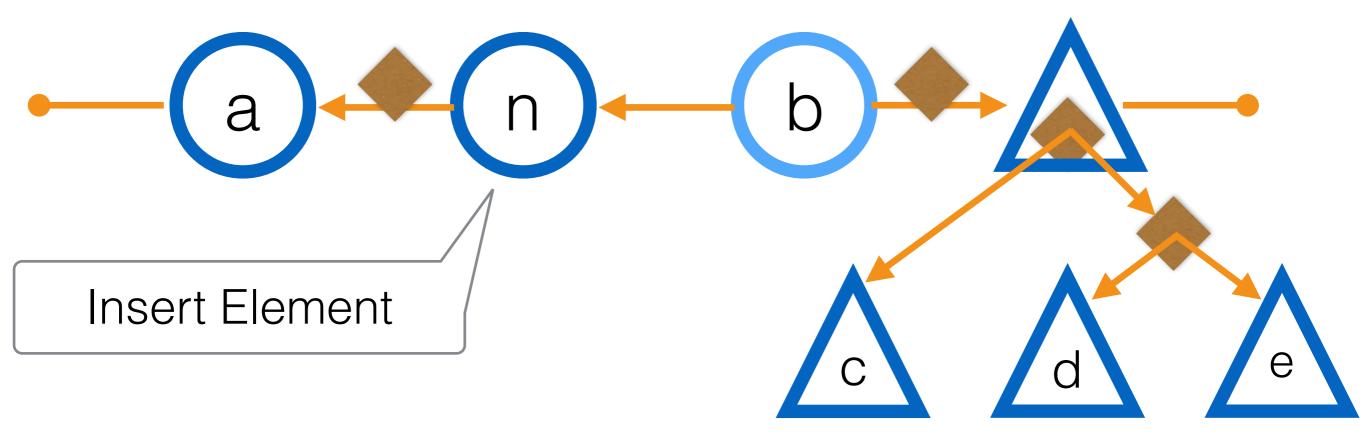


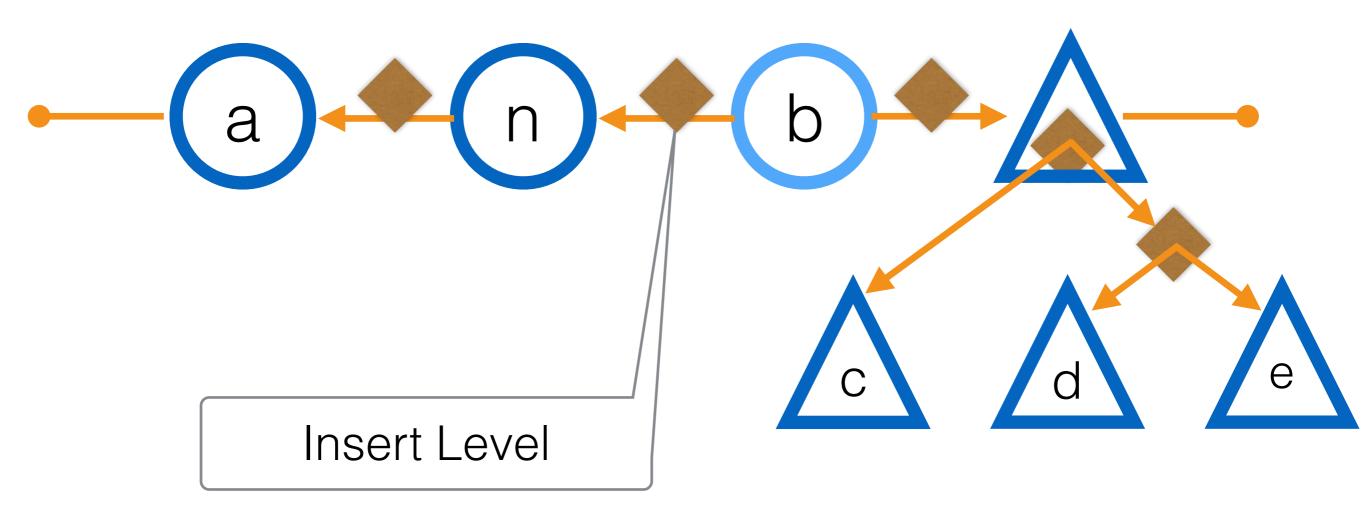
No Nil values in an unfocused RAZ

Closer look at our example

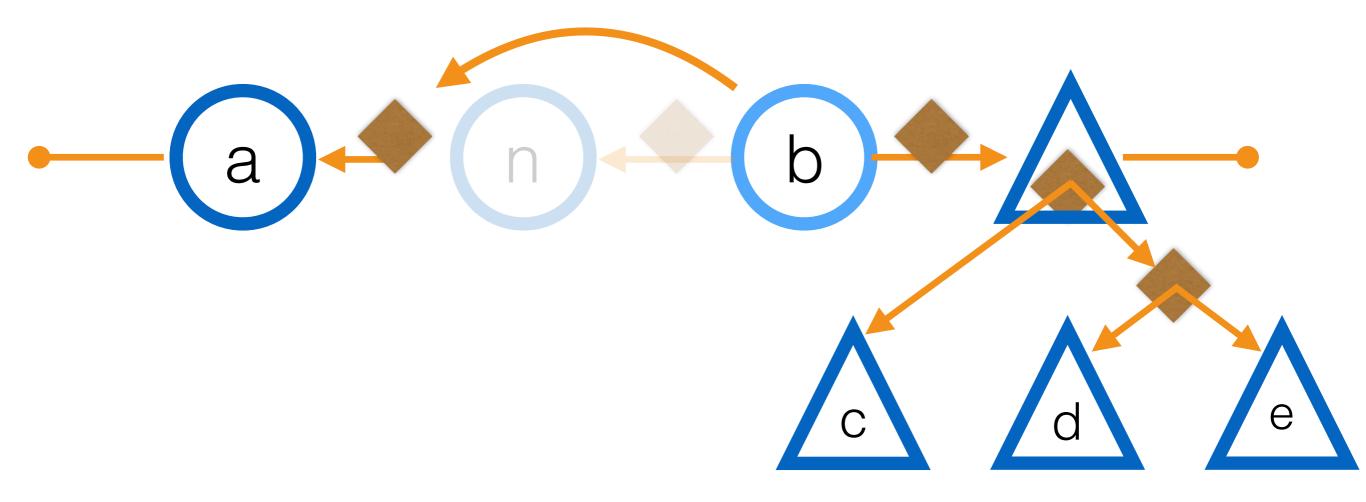


n а

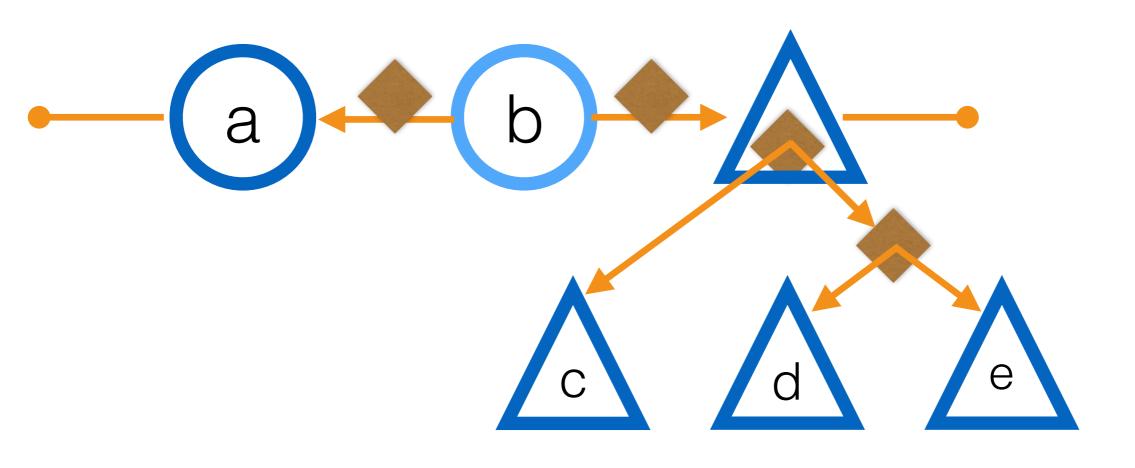


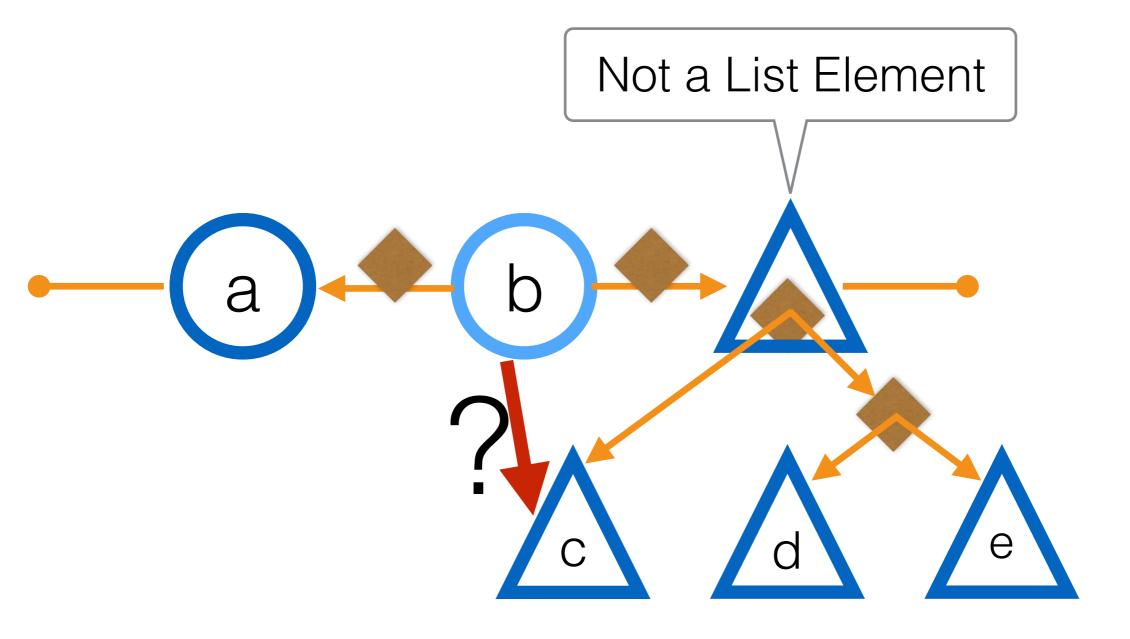


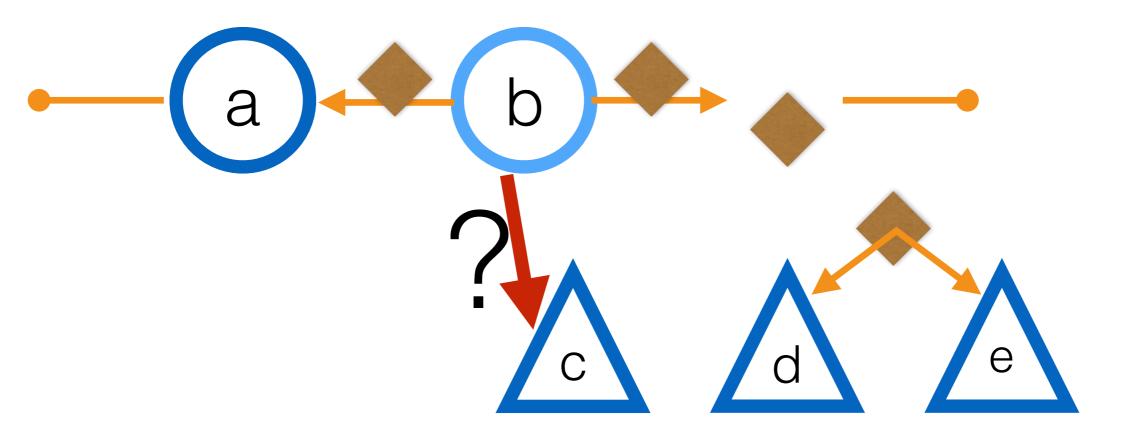
List-like Removal

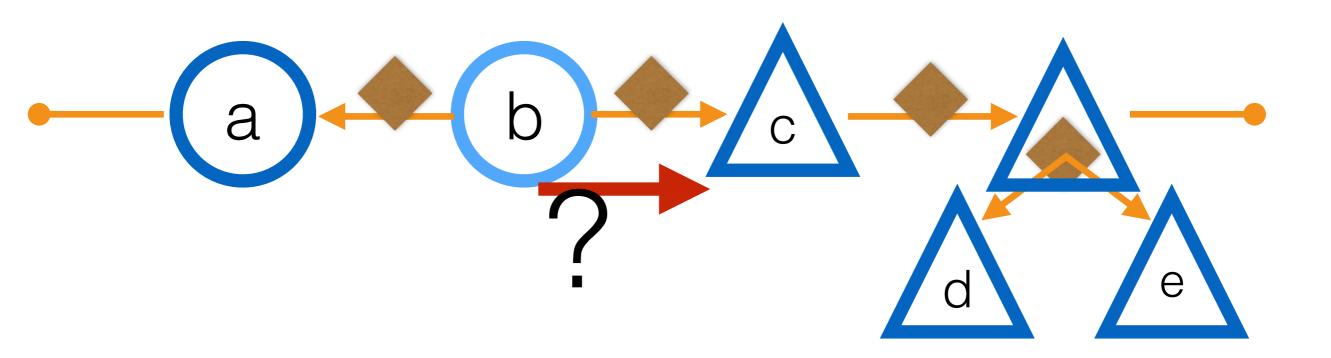


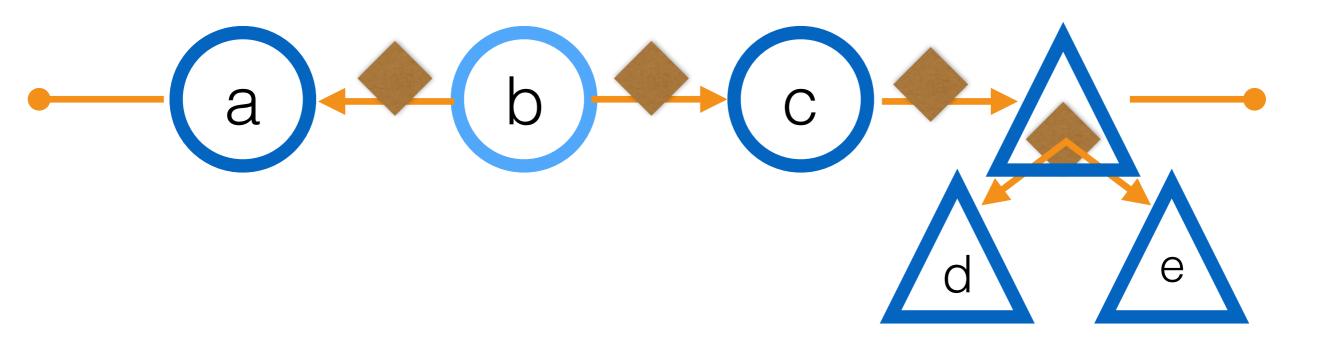
List-like Removal

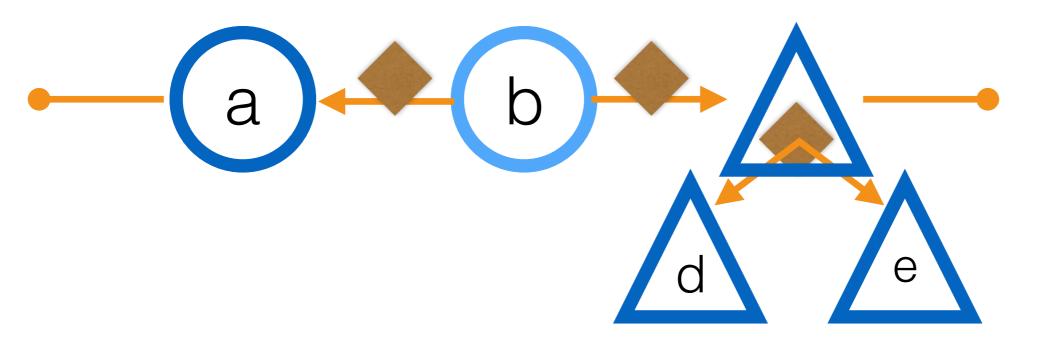


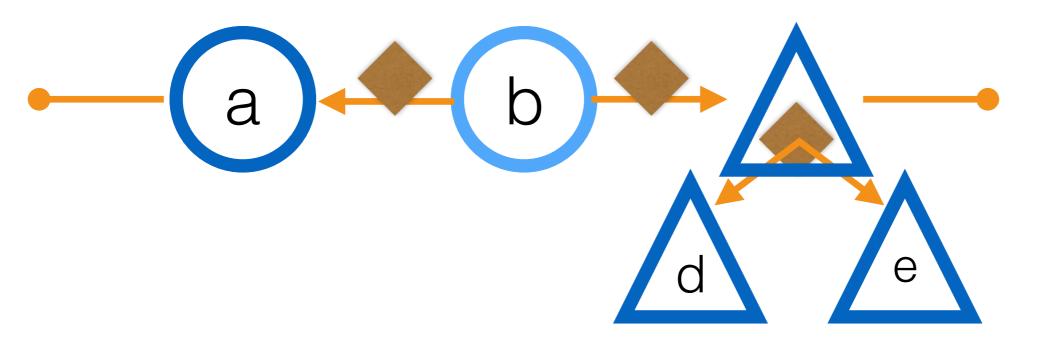


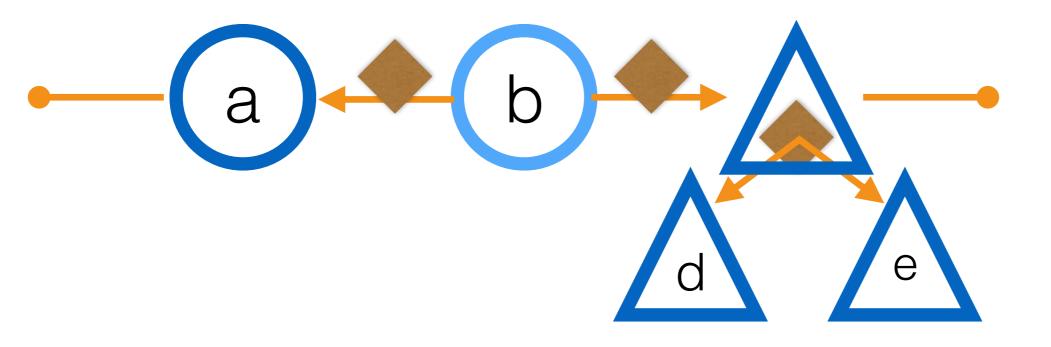


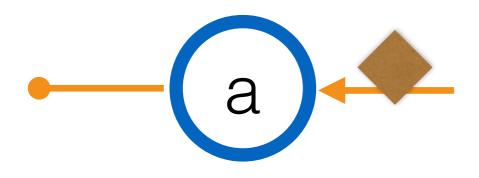


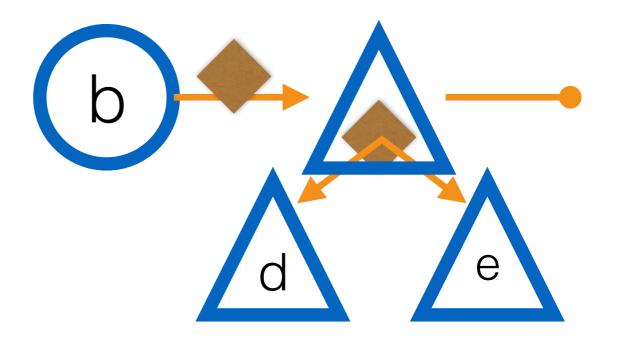


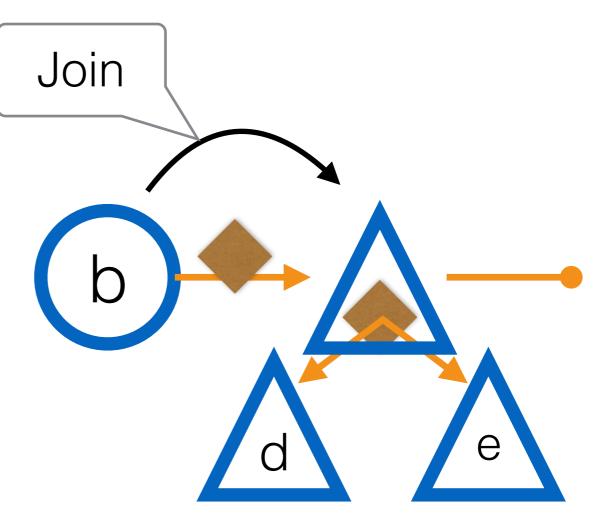


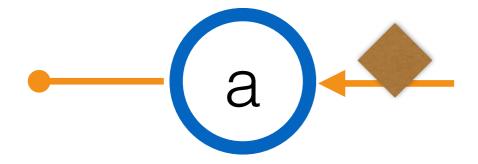


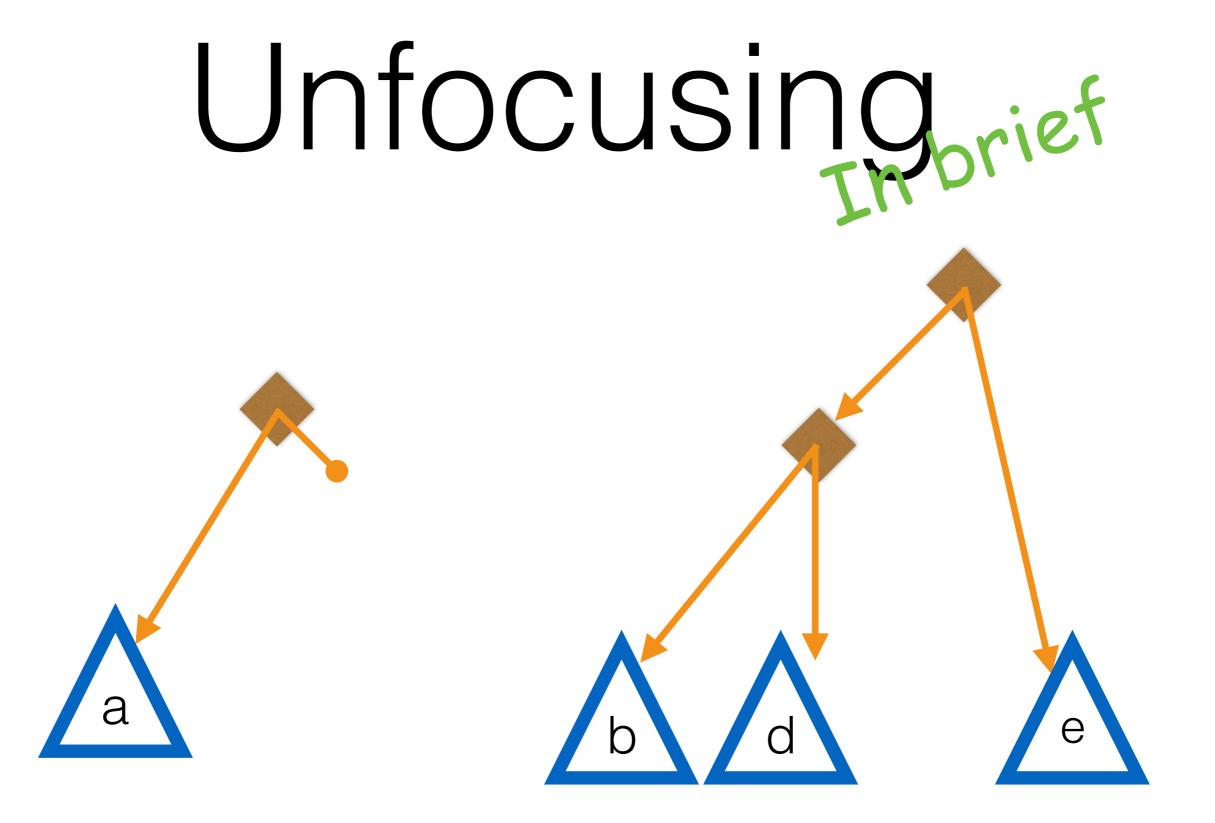


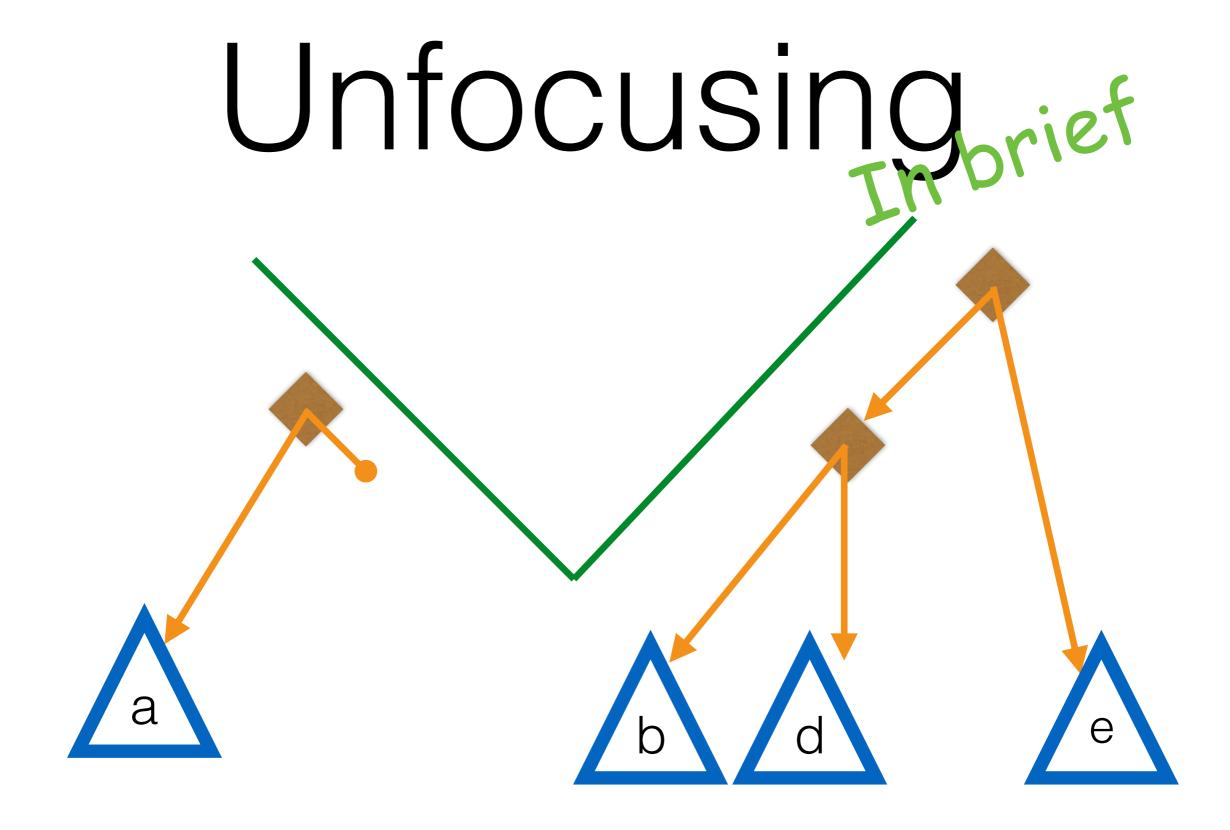


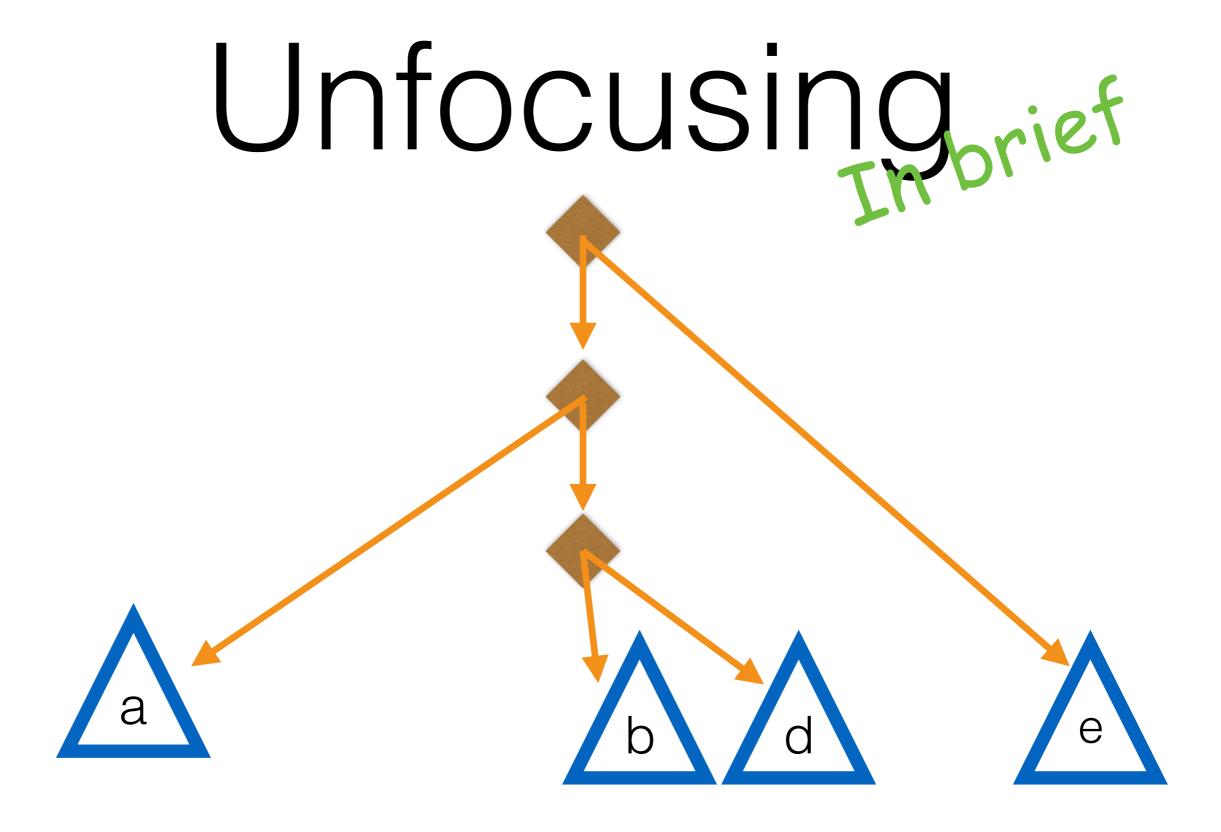


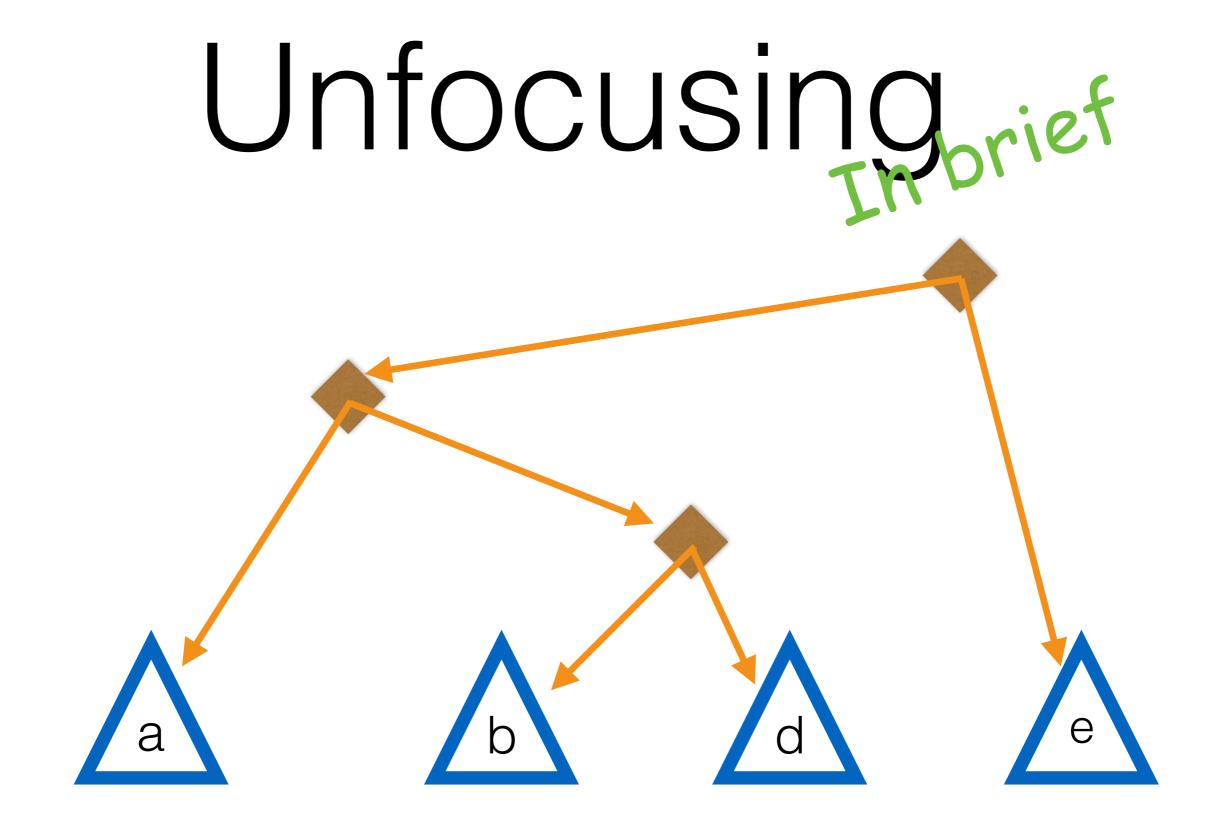


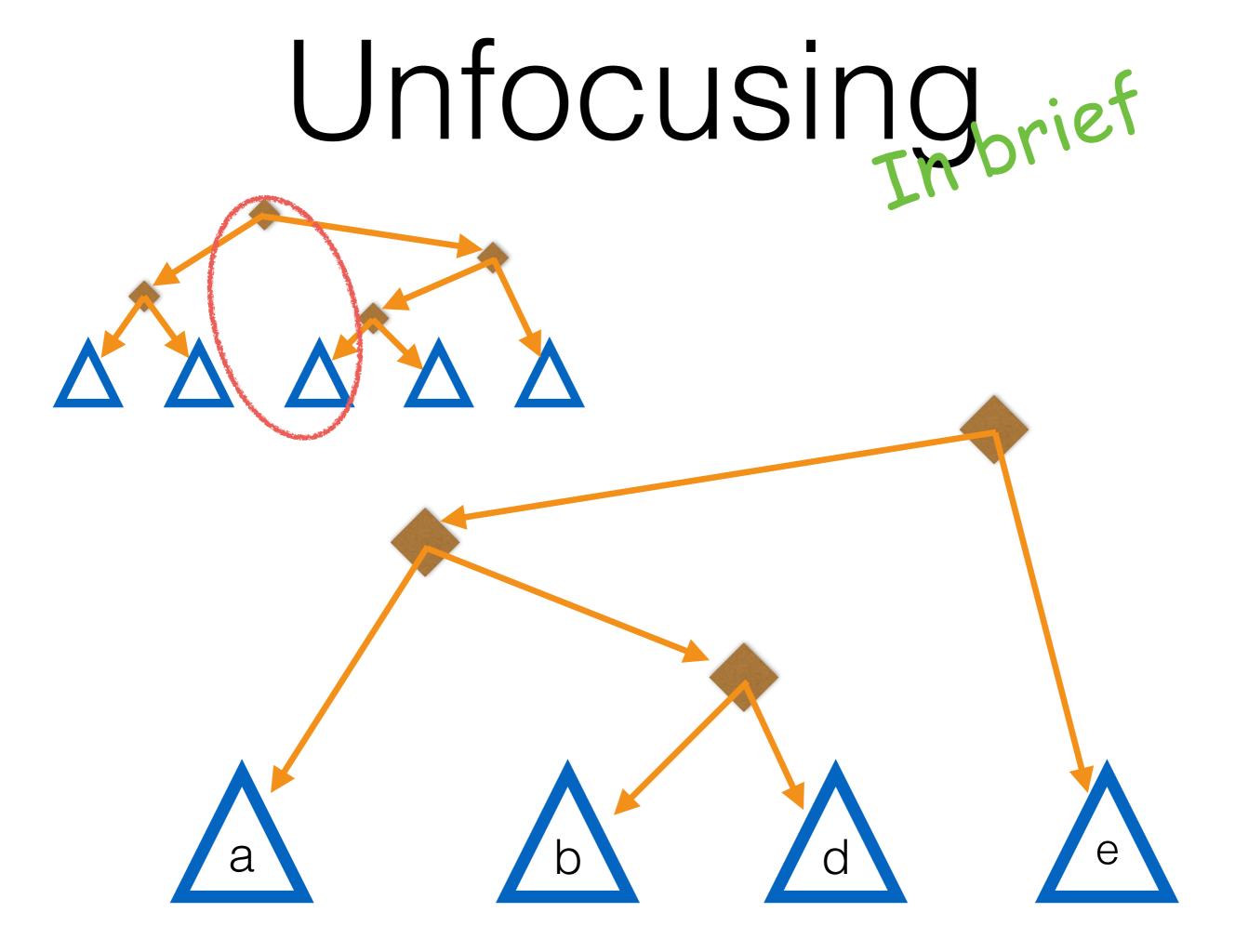












Switch to code

```
let focus pos tree =
  let rec focus = fun pos tree (l, r) \rightarrow
  match tree with
   Nil -> failwith "focus: internal Nil"
   Leaf(elm) ->
    assert (pos == 0);
    (l,elm,r)
  Bin(lv, _, branch_l, branch_r) ->
    let cnt = item_count branch l in
    if pos < cnt
    then
      focus pos branch_l
         (l, Level(lv, Tree(branch_r, r)))
    else
      let new_pos = (pos - cnt) in
      focus new_pos branch_r
         (Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

```
let focus pos tree =
  let rec focus = fun pos tree (l, r) \rightarrow
  match tree with
    Nil -> failwith "focus: internal Nil"
   Leaf(elm) ->
                                    Recursive function
    assert (pos == 0);
                                     with accumulator
    (l,elm,r)
   Bin(lv, _, branch_l, branch_f) ->
    let cnt = item count branch l in
    if pos < cnt
    then
      focus pos branch_l
         (l, Level(lv, Tree(branch_r, r)))
    else
      let new_pos = (pos - cnt) in
      focus new_pos branch_r
         (Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

let focus pos tree =
 let rec focus = fun pos tree (l, r) ->

match tree with

| Nil -> failwith "focus: internal Nil"

```
Leaf(elm) ->
assert (pos == 0);
(l,elm,r)
```

assert (pos == 0); A single pattern match!

Bin(lv, _, branch_l, branch_r) ->
let cnt = item_count branch_l in
if pos < cnt</pre>

then

focus pos branch_l

(l, Level(lv, Tree(branch_r, r)))
else

```
let new_pos = (pos - cnt) in
focus new_pos branch_r
```

```
(Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

```
let focus pos tree =
  let rec focus = fun pos tree (l, r) ->
  match tree with
    Nil -> failwith "focus: internal Nil"
   Leaf(elm) ->
    assert (pos == 0);
    (l,elm,r)
   Bin(lv, _, branch l, branch_r) ->
    let cnt = item_count branch l in
                         Return the accumulator as
    if pos < cnt
    then
                           a RAZ focused on this
      focus pos branch_l
         (l, Level(lv, Tree(branch r. element)
    else
      let new_pos = (pos - cnt) in
      focus new_pos branch_r
         (Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

```
let focus pos tree =
  let rec focus = fun pos tree (l, r) \rightarrow
  match tree with
    Nil -> failwith "focus: internal Nil"
   Leaf(elm) ->
    assert (pos == 0);
    (l,elm,r)
   Bin(lv, _, branch_l, branch_r) ->
    let cnt = item_count branch_l in
    if pos < cnt
                                     Branch based on
    then
                                    which side has the
      focus pos branch_l
         (l, Level(lv, Tree(branch r,focus element
    else
      let new_pos = (pos - cnt) in
      focus new_pos branch_r
         (Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

```
let focus pos tree =
  let rec focus = fun pos tree (l, r) ->
  match tree with
    Nil -> failwith "focus: internal Nil"
   Leaf(elm) ->
                                     On left, recurse on
    assert (pos == 0);
    (l,elm,r)
                                    left branch with right
  Bin(lv, _, branch_l, branch_r) -> branch in right
    let cnt = item count branch l
                                   in
                                       accumulator
    if pos < cnt
    then
      focus pos branch l
         (l, Level(lv, Tree(branch_r, r)))
    else
      let new_pos = (pos - cnt) in
      focus new_pos branch_r
         (Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

```
let focus pos tree =
  let rec focus = fun pos tree (l, r) \rightarrow
  match tree with
    Nil -> failwith "focus: internal Nil"
    Leaf(elm) ->
    assert (pos == 0);
    (l,elm,r)
   Bin(lv, _, branch_l, branch_r) ->
    let cnt = item count branch l in
                                     On right, recurse on
    if pos < cnt
    then
                                     right branch with left
      focus pos branch_l
                                        branch in left
         (l, Level(lv, Tree(branch
                                     r,
                                        accumulator
    else
      let new_pos = (pos - cnt) in/
      focus new_pos branch_r
         (Level(lv, Tree(branch_l, l)), r)
in focus pos tree (Nil, Nil)
```

let alter : dir -> 'a -> 'a raz -> 'a raz =
 let rec alter new side zip = match zip with
 | Nil -> failwith "alter: past end of seq"
 | Cons(_,rest) -> Cons(new,rest)
 | Level(lv,rest) -> Level(lv,alter new side rest)
 | Tree _ -> alter new side (trim side zip)
 in fun side elm (l,e,r) -> match side with
 | L -> (alter elm L l,e,r)
 | R -> (l,e,alter elm R r)

Local edits take directions

let alter : dir -> 'a -> 'a raz -> 'a raz =

let rec alter new side zip = match zip with
| Nil -> failwith "alter: past end of seq"
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| L -> (alter elm L l,e,r)
| R -> (l,e,alter elm R r)

let alter : dir -> 'a -> 'a raz -> 'a raz = let rec alter new side zip = match zip with | Nil -> failwith "alter: past end of seq" | Cons(_,rest) -> Cons(new,rest) | Level(lv,rest) -> Level(lv,alter new side rest) | Tree _ -> alter new side (trim side zip) in fun side elm (l,e,r) -> match side with | L -> (alter elm L l,e,r) | R -> (l,e,alter elm R r) |

Two pattern matches

let alter : dir -> 'a -> 'a raz -> 'a raz = let rec alter new side zip = match zip with Nil -> failwith "alter: past end of seq" Cons(,rest) -> Cons(new,rest) Level(lv,rest) -> Level(lv,alter new side rest) | Tree _ -> alter new side (trim side zip) in fun side elm (l,e,r) -> match side with | L -> (alter elm L l,e,r)| R -> (l,e,alter elm R r)Edit is similar for each side

let alter : dir -> 'a -> 'a raz -> 'a raz =
 let rec alter new side zip = match zip with
 | Nil -> failwith "alter: past end of seq"
 Cons(_,rest) -> Cons(new,rest)
 Level(lv,rest) -> Level(lv,alter new side rest)
 I Tree _ -> alter new side (trim side zip)
 in fun side elm (l,e,r) -> match side with
 L -> (alter elm L l,e,r)
 R -> (l,e,alter elm R r)

Two common cases

We always reach a level first

let alter : dir -> 'a -> 'a raz -> 'a raz =
 let rec alter new side zip = match zip with
 | Nil -> failwith "alter: past end of seq"
 | Cons(_,rest) -> Cons(new,rest)
 | Level(lv,rest) -> Level(lv,alter new side rest)
 | Tree _ -> alter new side (trim side zip)
 in fun side elm (l,e,r) -> match side with
 | L -> (alter elm L l,e,r)
 | R -> (l,e,alter elm R r)

let alter : dir -> 'a -> 'a raz -> 'a raz = let rec alter new side zip = match zip with Nil -> failwith "alter: past end of seq" Cons(_,rest) -> Cons(new,rest) Level(lv,rest) -> Leve (lv,alter new side rest) Tree _ -> alter new side (trim side zip) in fun side elm $(l,e,r) \rightarrow$ match side with | L -> (alter elm L l,e,r)| R -> (l,e,alter elm R r)

Return new element

RAZ in OCaml

RAZ in OCaml

Fingertree in OCaml

RAZ in OCaml

Insertion and removal at random point

Fingertree in OCaml

RAZ in OCaml

Insertion and removal at random point

Fingertree in OCaml

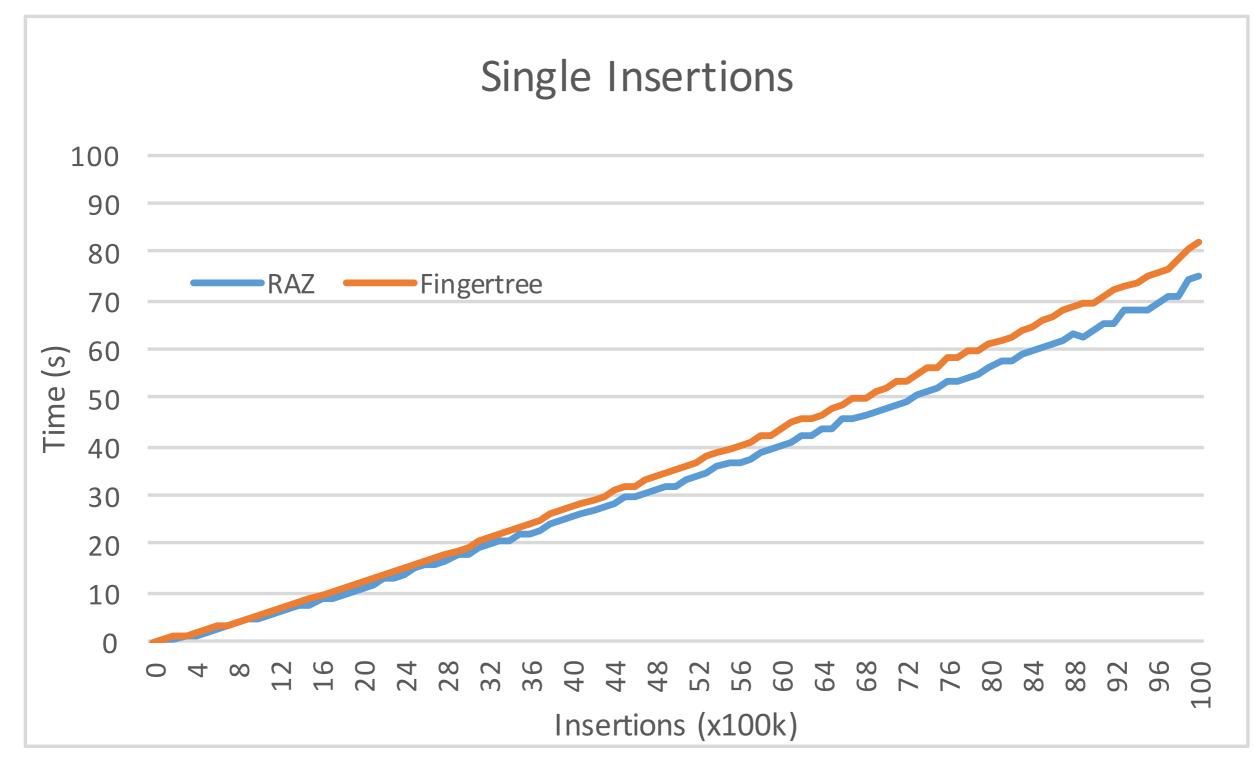
Insertion at random point

Insertion, Removal

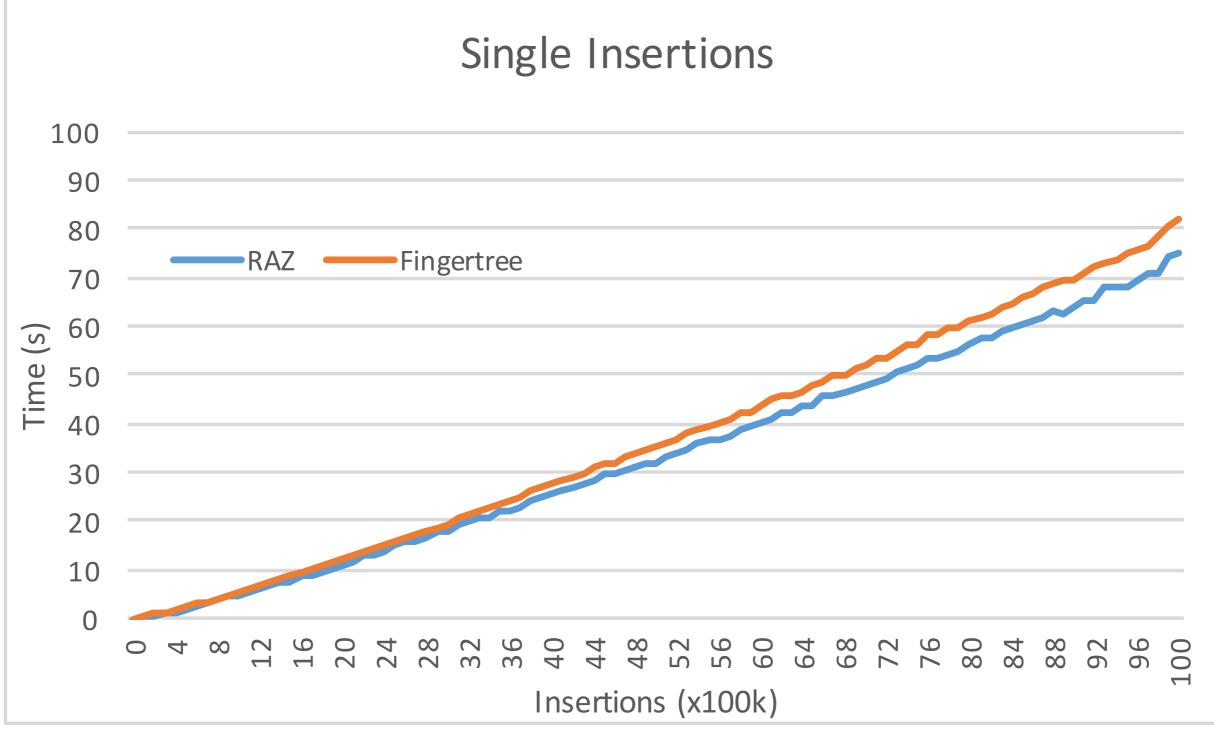
Insertions and Deletions



Insertion at random



Insertion at random



Simplicity as performance?

Random Access Zipper

- Accessible
- Editable
- Simple
- Fast

Random Access Zipper

Simple enough to include these principles in your own data types!