

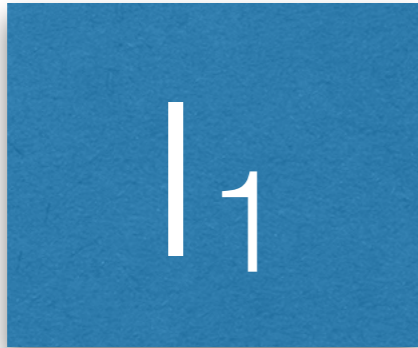
IODyn: A High-Level Language for Incremental Computation

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1 kyleheadley.github.io

What is Incremental Computation?

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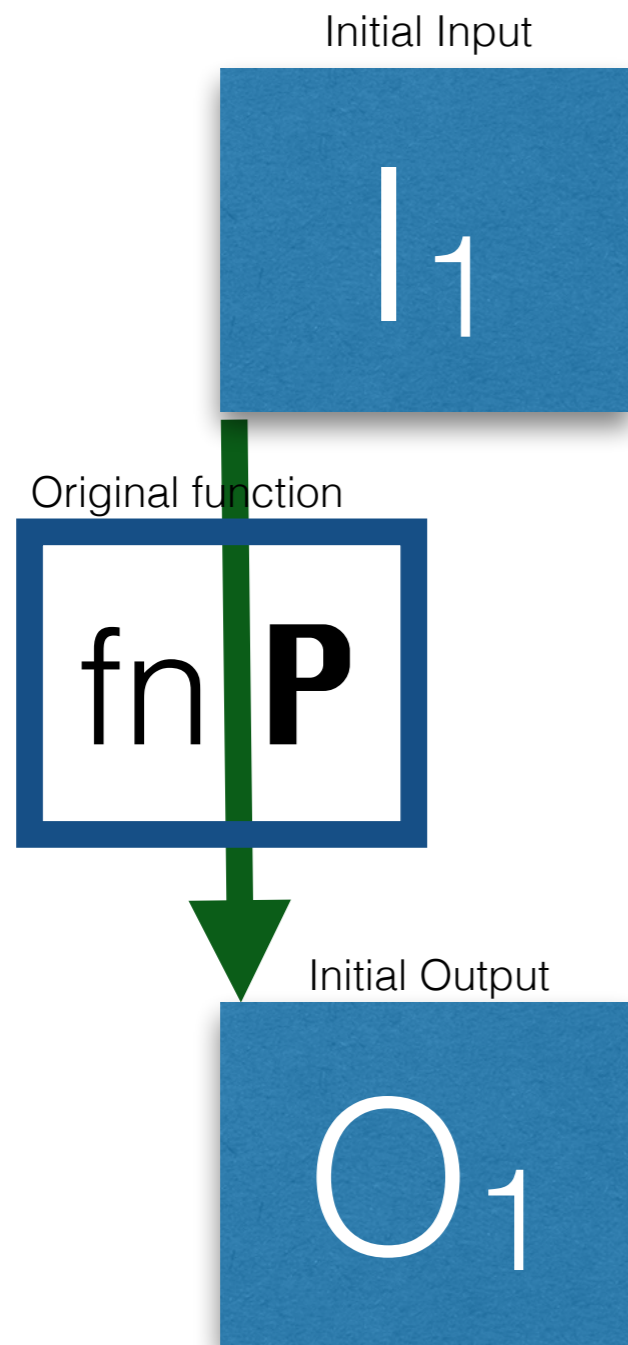
Initial Input



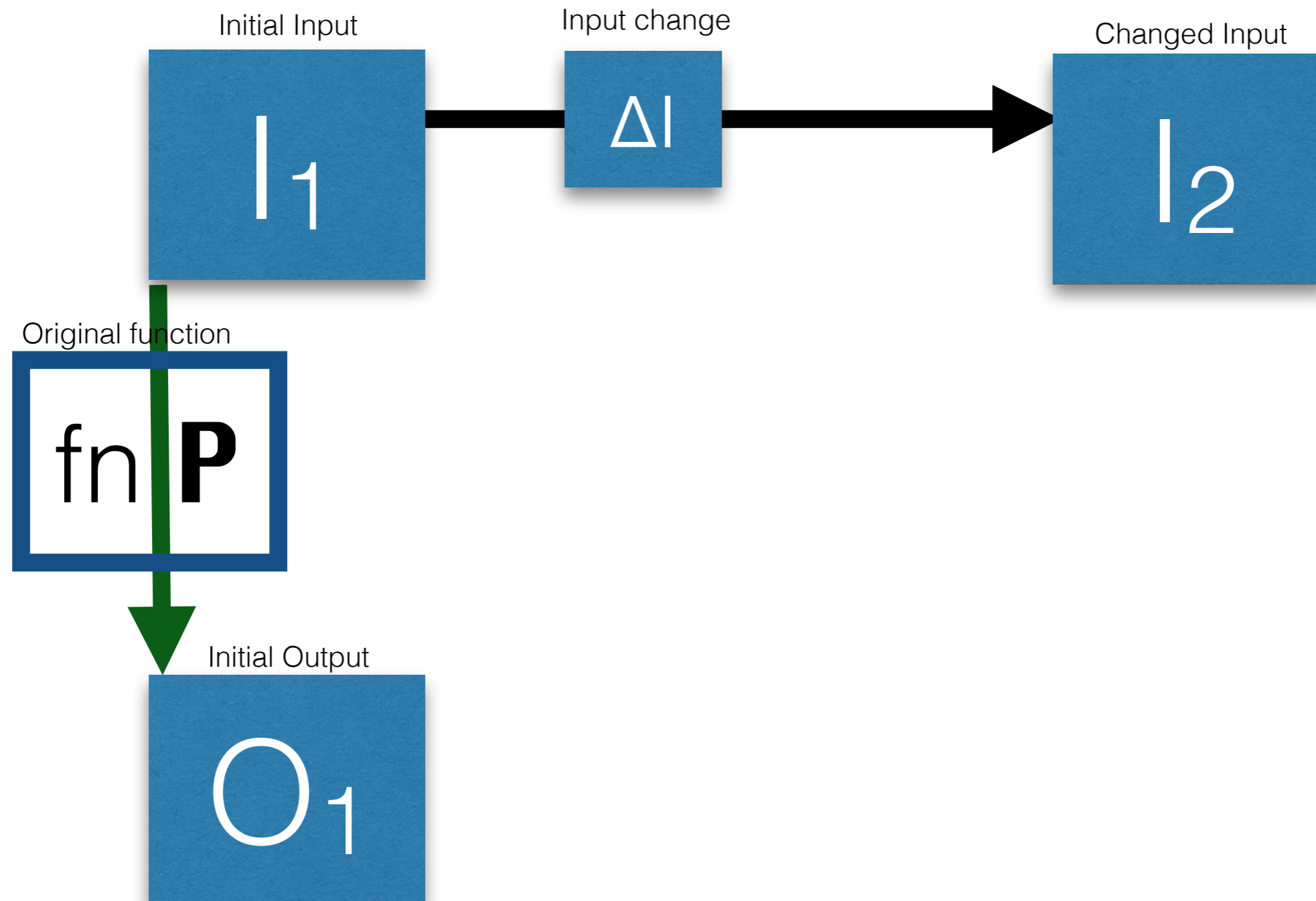
Original function



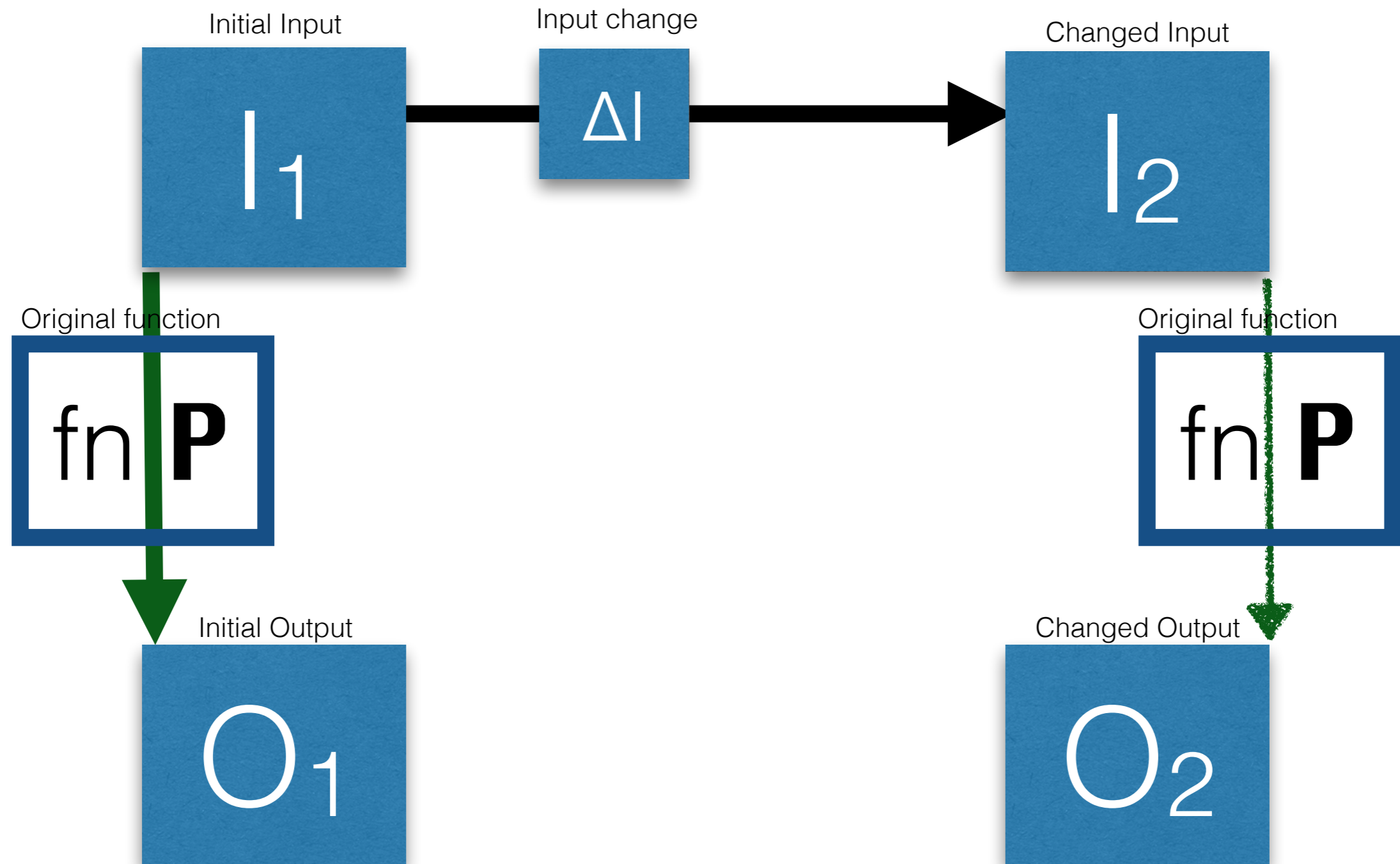
What is Incremental Computation?



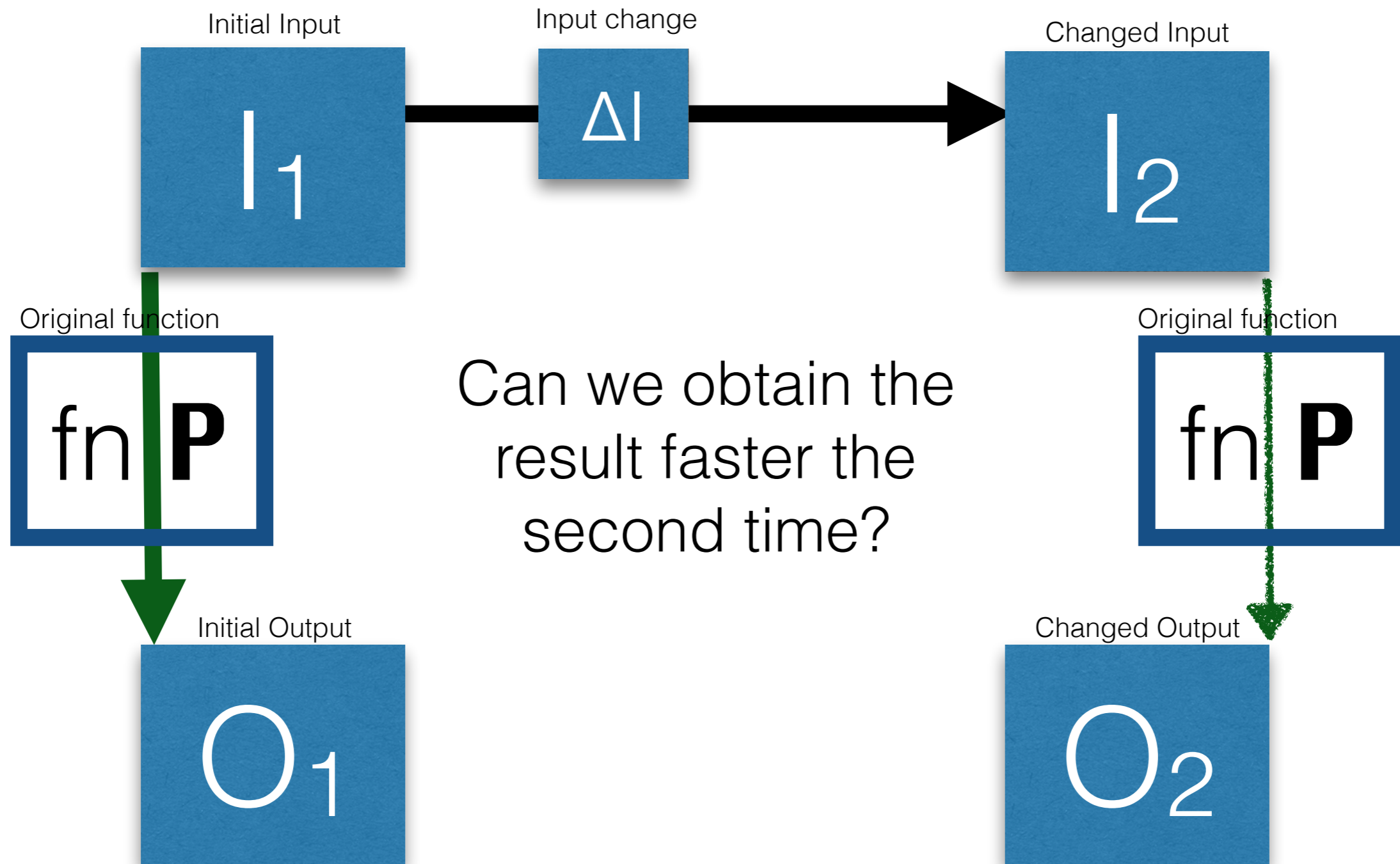
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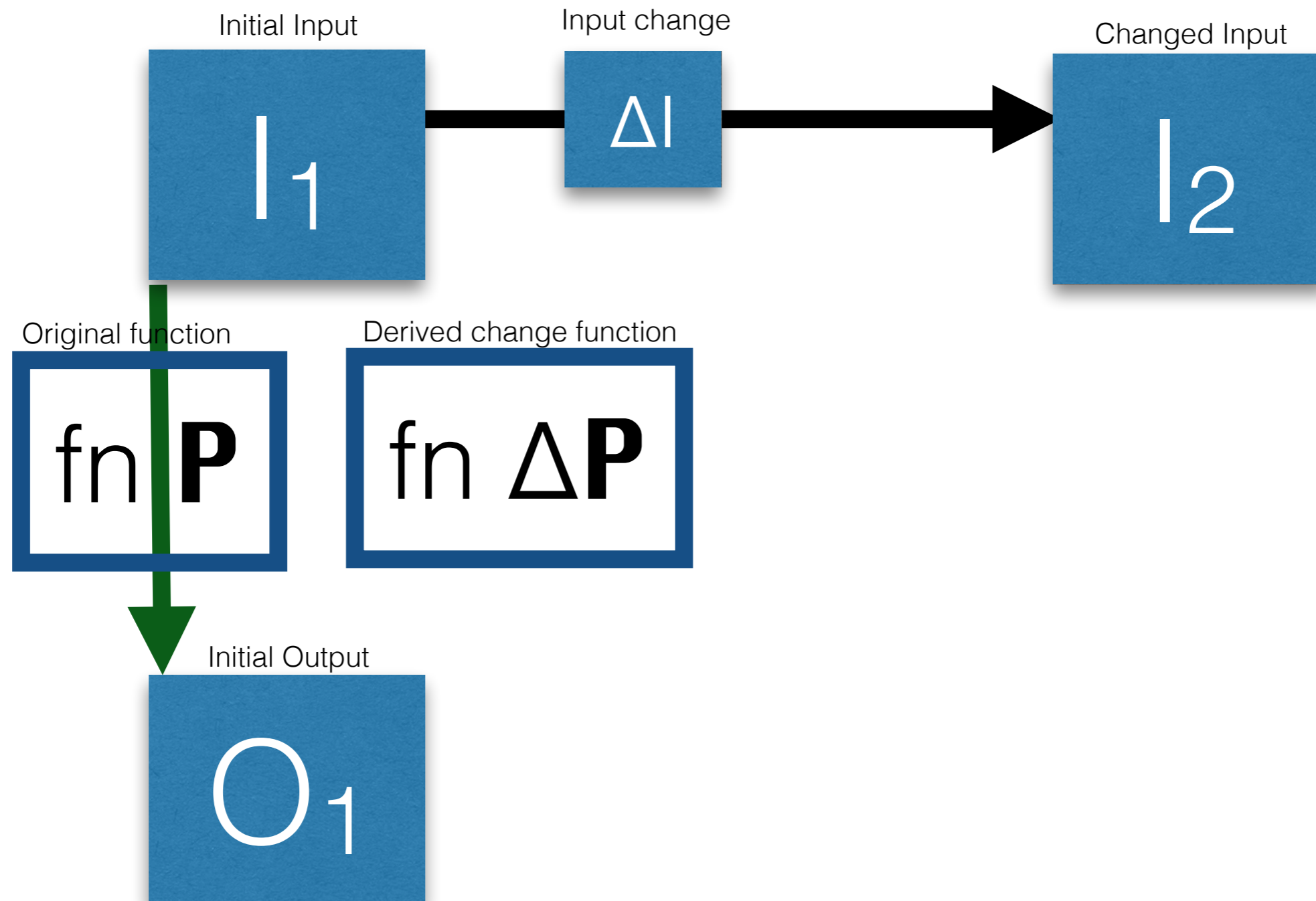
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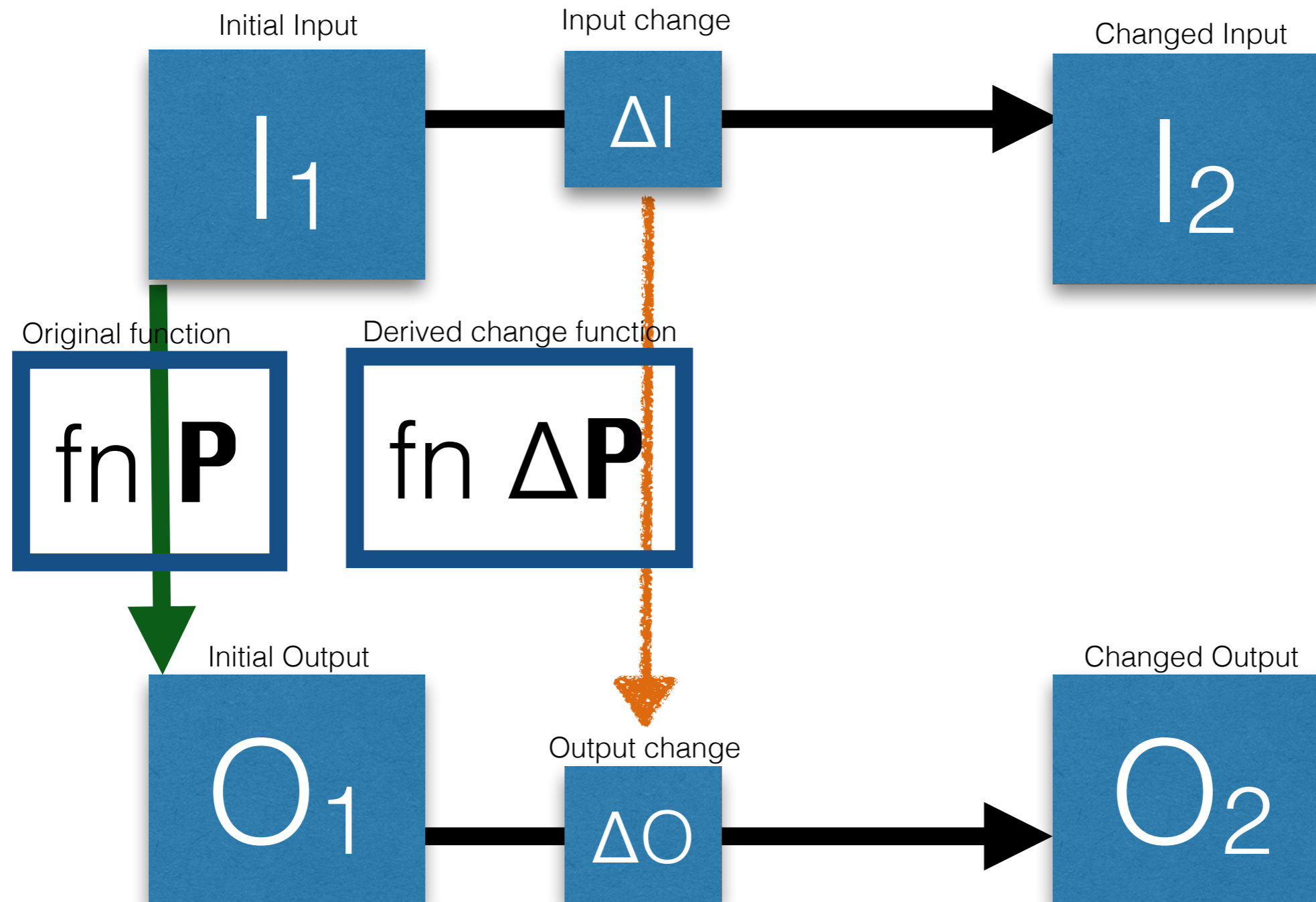
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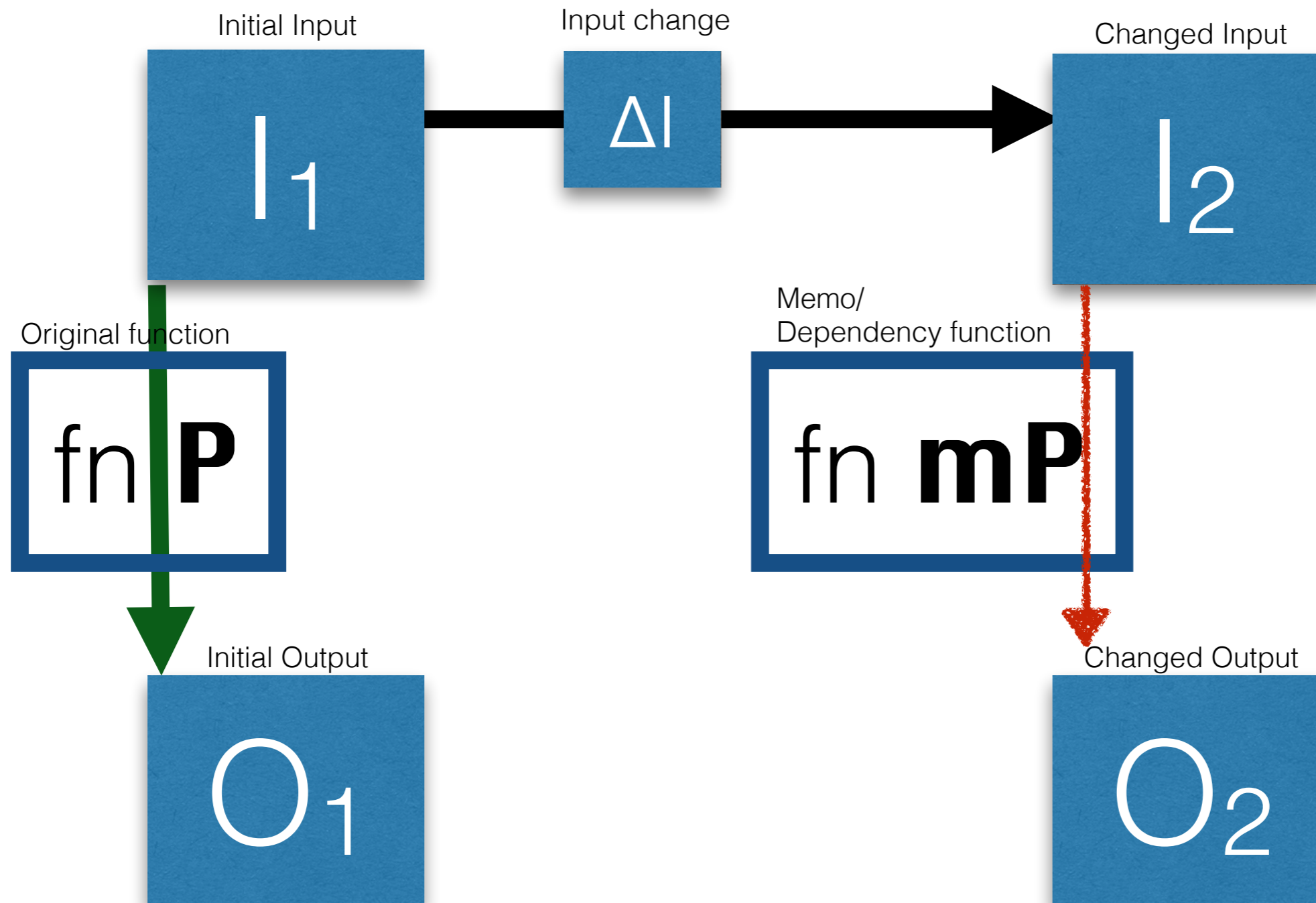
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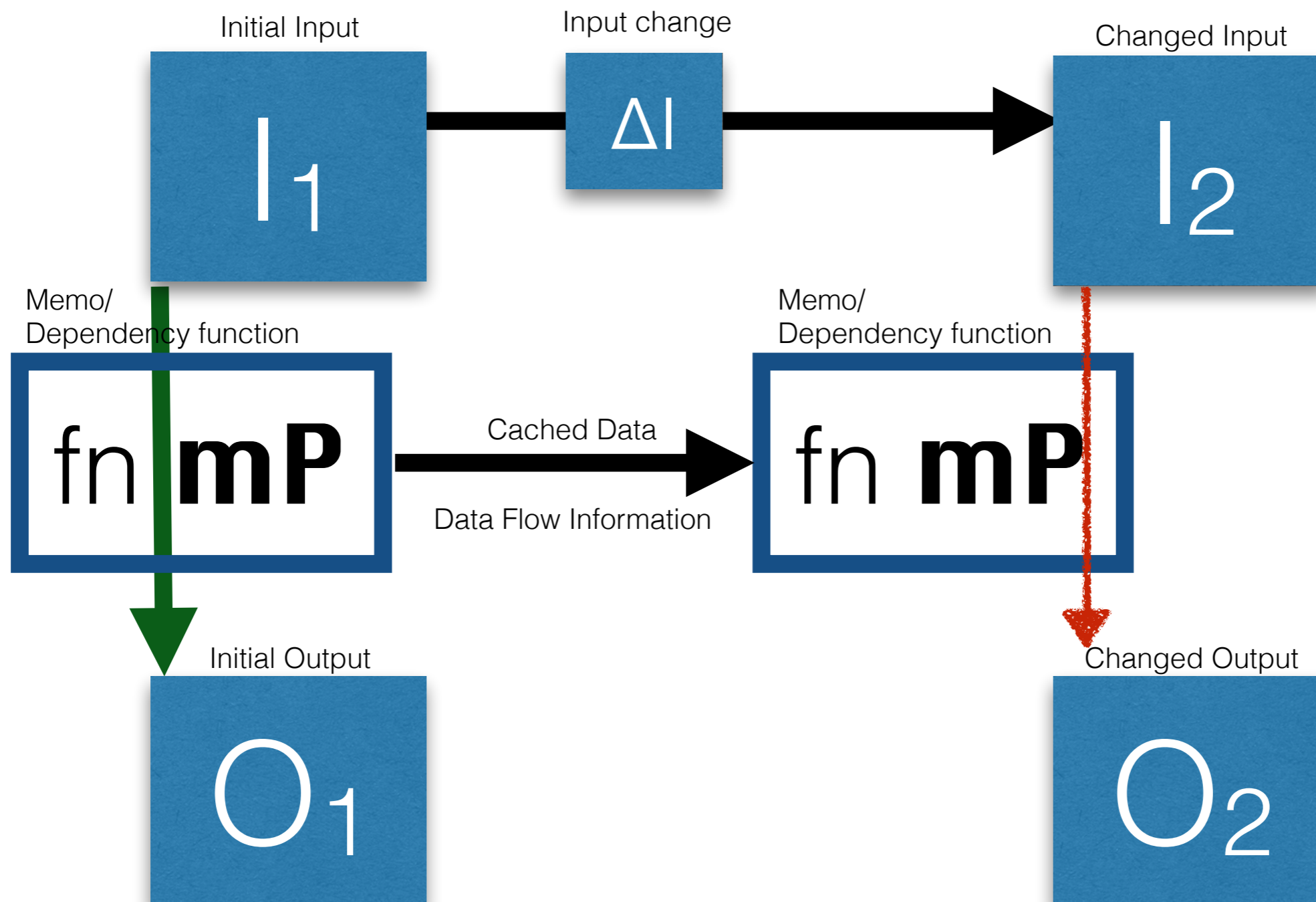
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How do we use Incremental Computation?

General-purpose incremental computation engines are accessed through libraries

Programmers annotate variables and functions based on their knowledge of the program and its incremental behavior

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But library use may not be straightforward and improper
caching can lead to slowdowns or incorrect results

How do we improve Incremental Computation?

Speedups

Correctness

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Speedups

Caching strategies that work efficiently for common computation tasks

Correctness

Static guarantees of proper use of incremental features

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Static guarantees of proper use of incremental features

Our lab is working on:

Incremental
Collections
Libraries

Type system for
use of library
features

How do we improve Incremental Computation?

Speedups

Caching strategies that work efficiently for common computation tasks

Correctness

Static guarantees of proper use of incremental features

Our lab is working on:

Incremental
Collections
Libraries

Type system for
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But these are specific solutions, with faults:

Conflicts between multiple applied incremental functions may still interfere with cache behavior

Writing the code is still difficult, and it provides no guarantees of performance

IODyn

Implicitly Incremental Language

Work in progress:

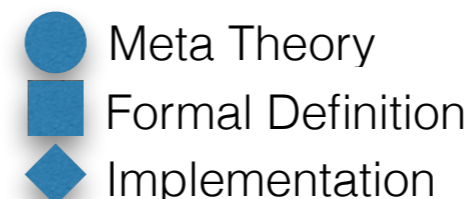
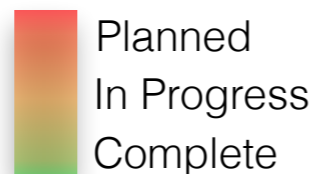
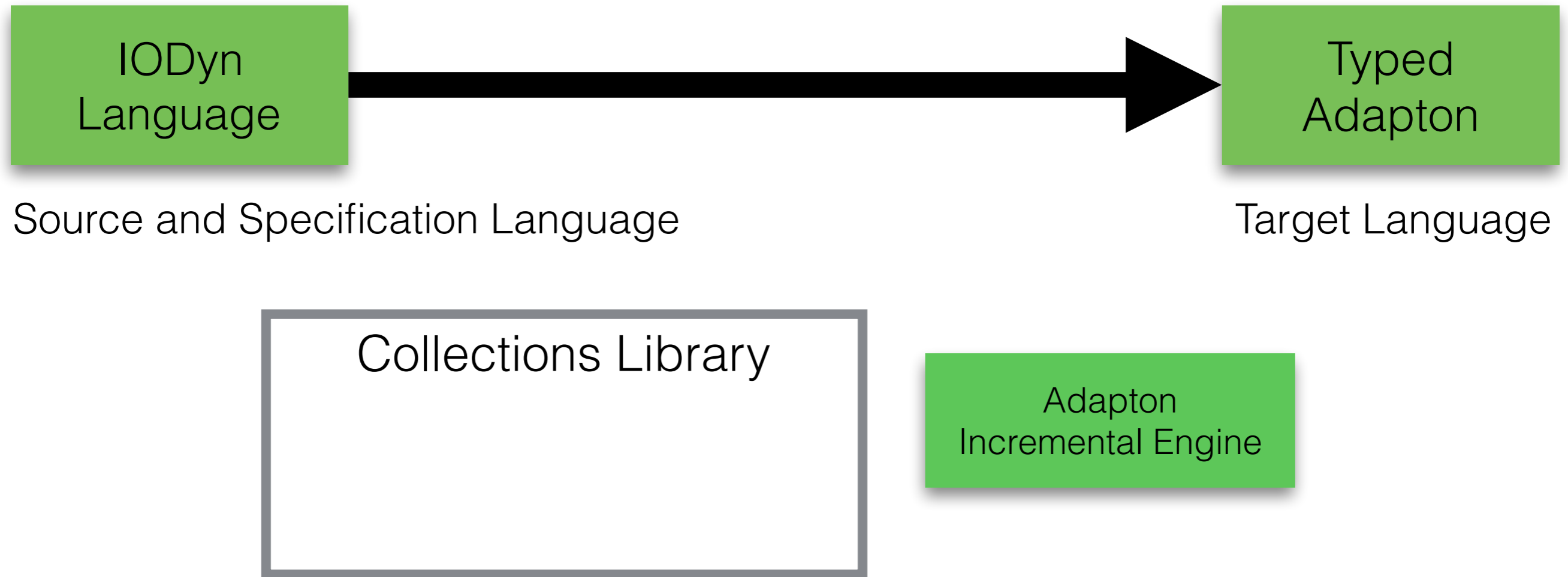
github.com/cuplv/iodyn-lang.rust

Based on the lambda-calculus

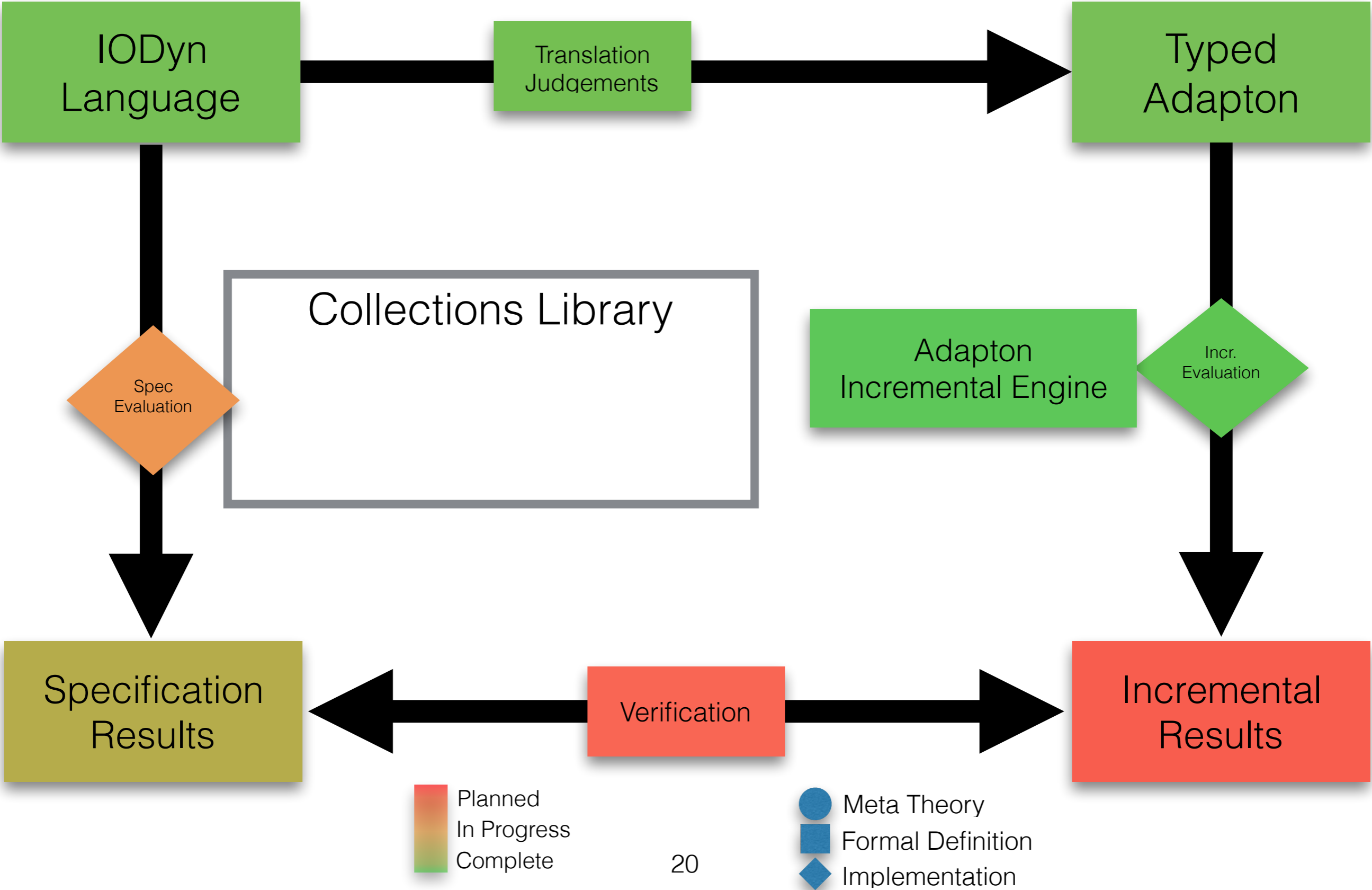
Collections primitives

Translated to Typed Adaption

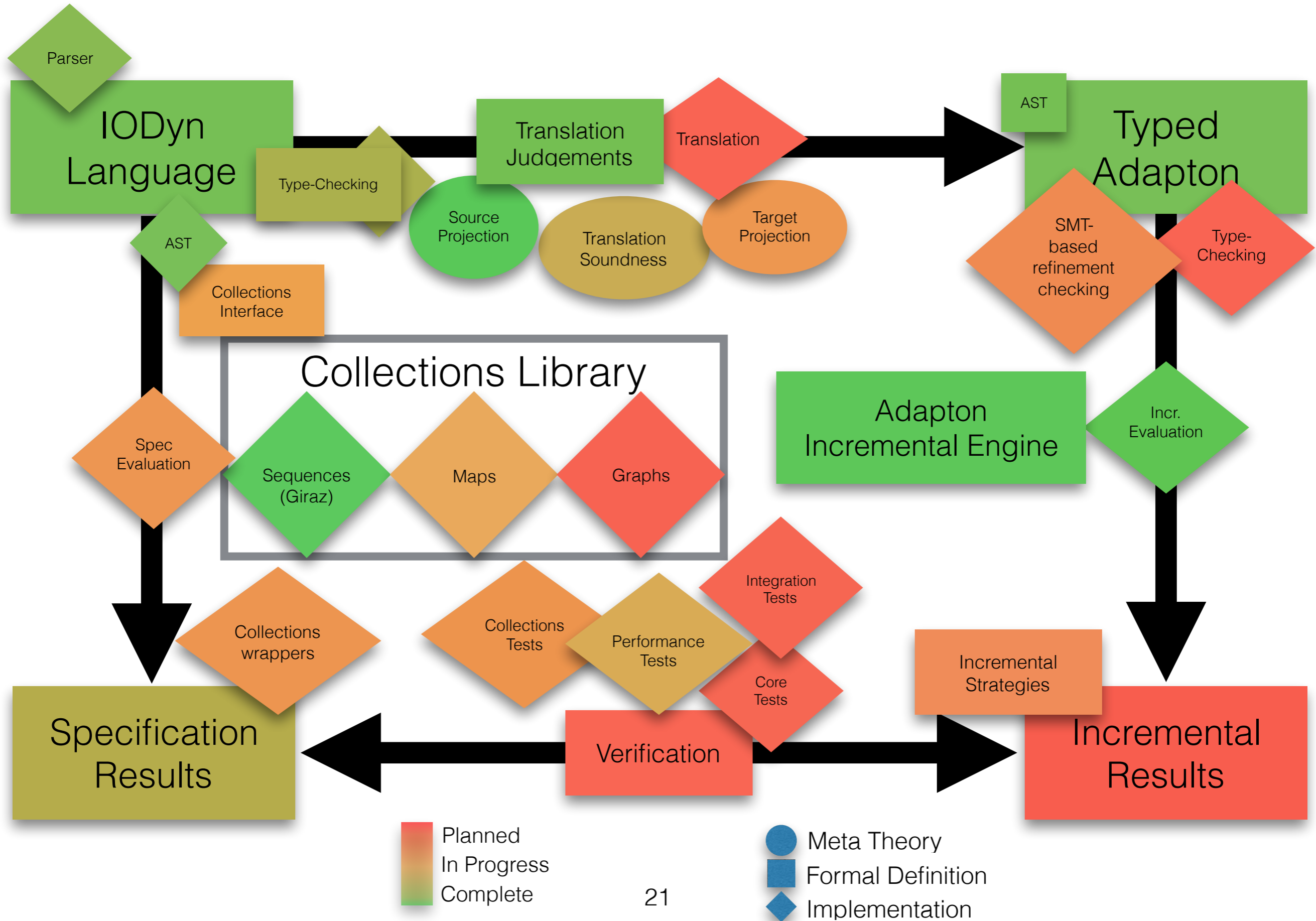
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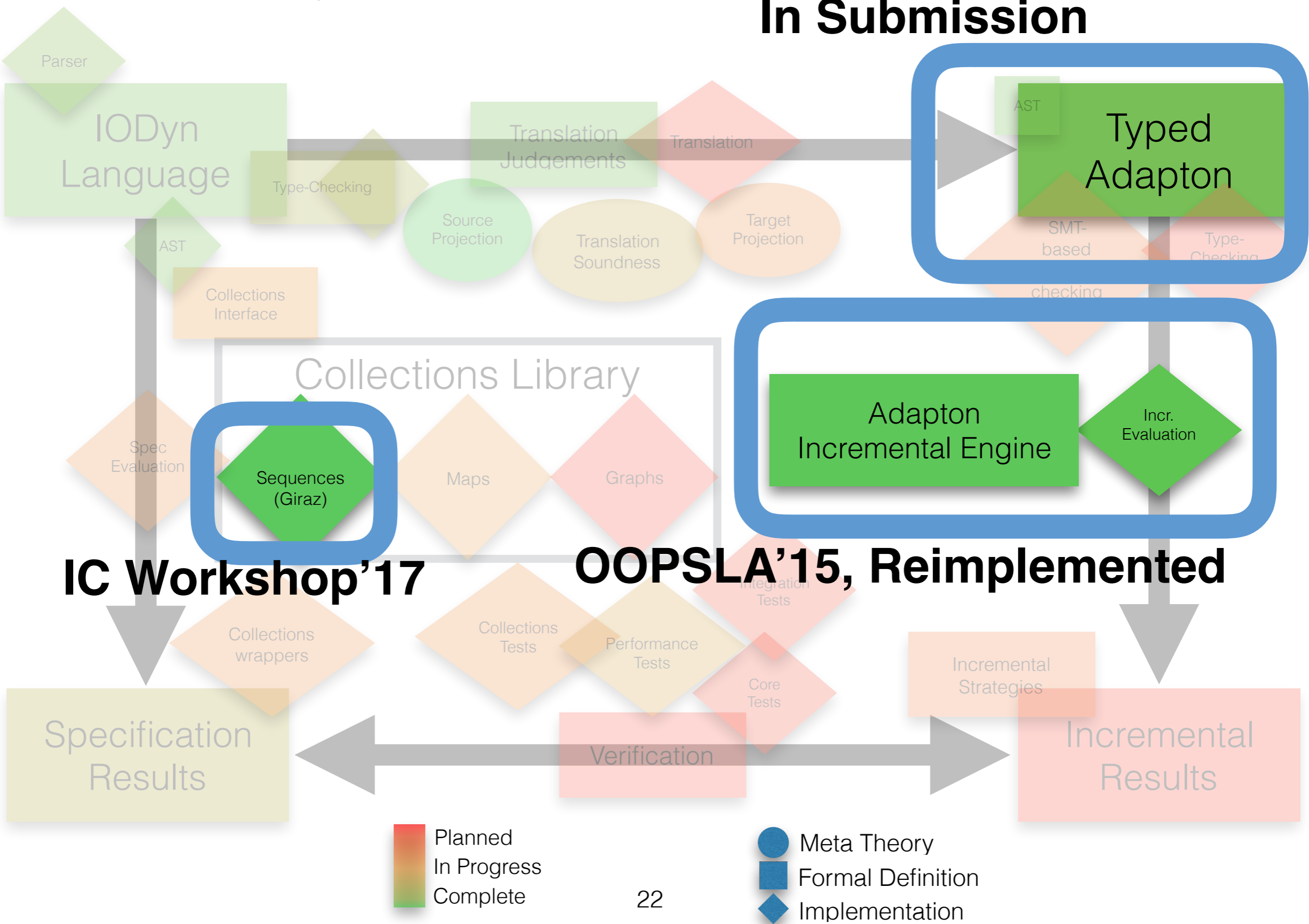
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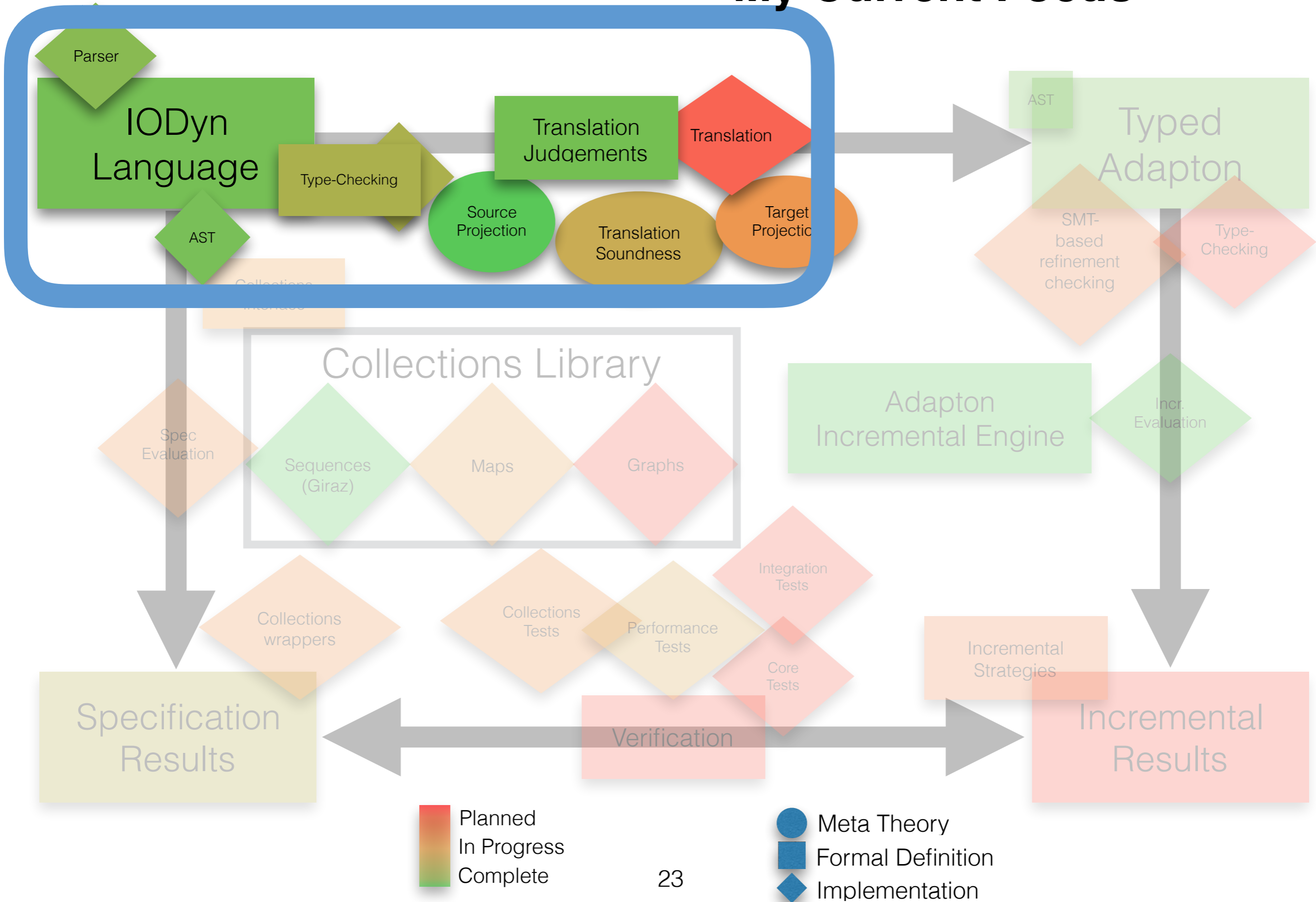


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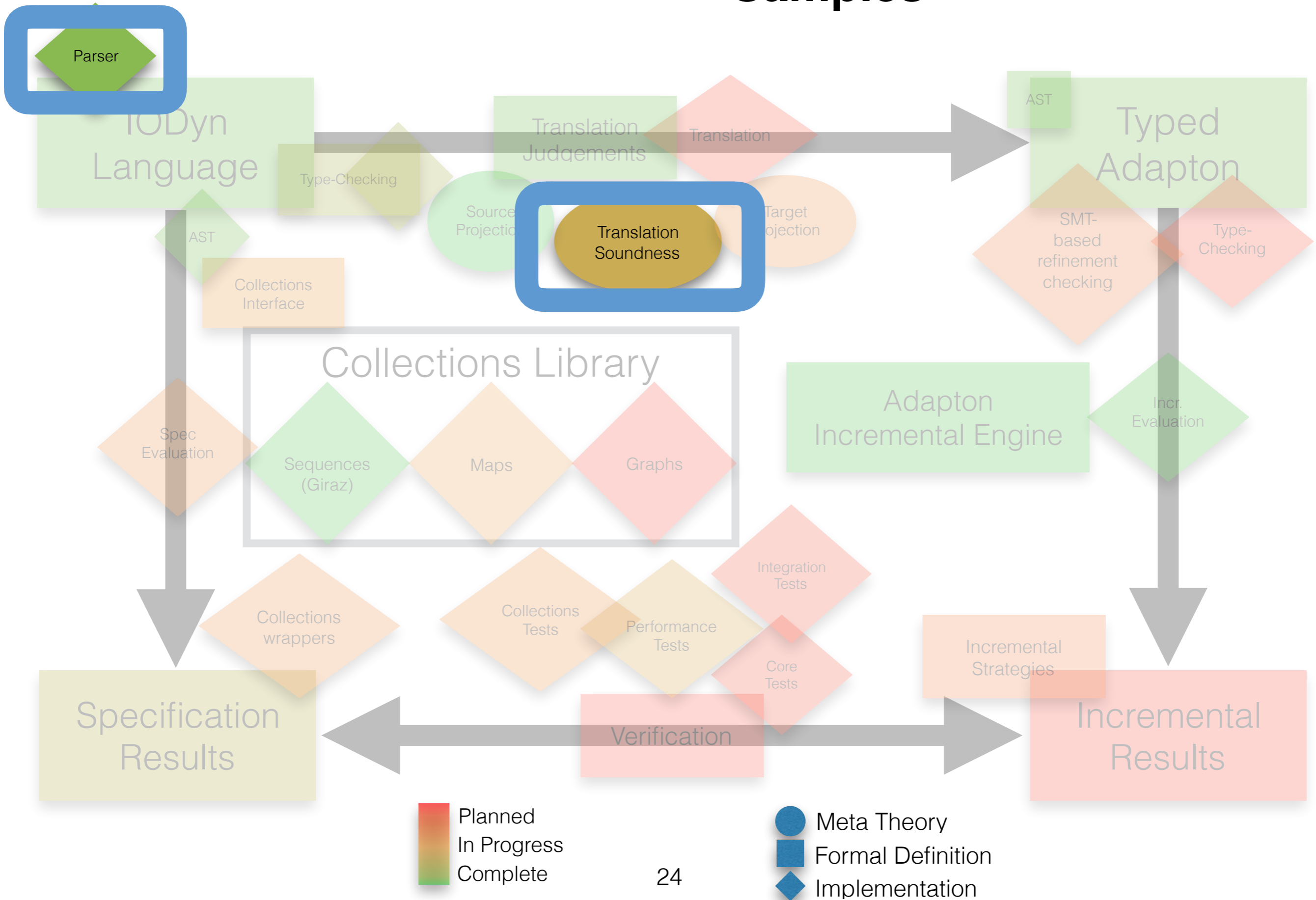
What is IODyn?

My Current Focus



What is IODyn?

Samples



Parsing with Rust Macros

Parsing with Rust Macros

Defined by a set of rewriting rules

```
macro_rules! macro_name {  
    { pattern } => { replacement };  
    ...  
}  
  
macro_name! [ pattern ]  
// replacement
```

Parsing with Rust Macros

Defined by a set of rewriting rules

```
macro_rules! macro_name {  
    { literal1 literal2 } => { literal3 };  
    { literal1 $variable:tt } => { $variable literal1 };  
    { literal2 $($repeats:tt),+ } => { $(literal2 $repeats);+ }  
}  
  
macro_name! [ literal1 something ]  
// something literal1  
macro_name! [literal2 r, s, t]  
// literal2 r; literal2 s; literal2 t
```

Parsing with Rust Macros

Sample rule

```
macro_rules! make_exp {  
  // lam r.e (lambda)  
  {lam $var:ident . $($body:tt)+ } =>  
  {{Exp::Lam(  
    stringify![$var].to_string(),  
    Rc::new(make_exp![$($body)+])  
  )  
  ...  
  }};  
}
```

Sample code

```
{fix qh. lam pts. lam line. lam hull.  
let complete = { SeqIsEmpty(pts) }  
if (complete) then { ret hull } else {  
  ...  
}} : Seq((nat x nat)) ->  
  ((nat x nat) x (nat x nat)) ->  
  Seq((nat x nat)) ->  
  F Seq((nat x nat))
```

Parsing with Rust Macros

$A \rightarrow B \rightarrow C \rightarrow D$

Not directly possible, instead,
use a fold, checking each token

Parsing with Rust Macros

A -> B -> C -> D

⇓ Fold, find `->`

(A) (B) (C) (D)

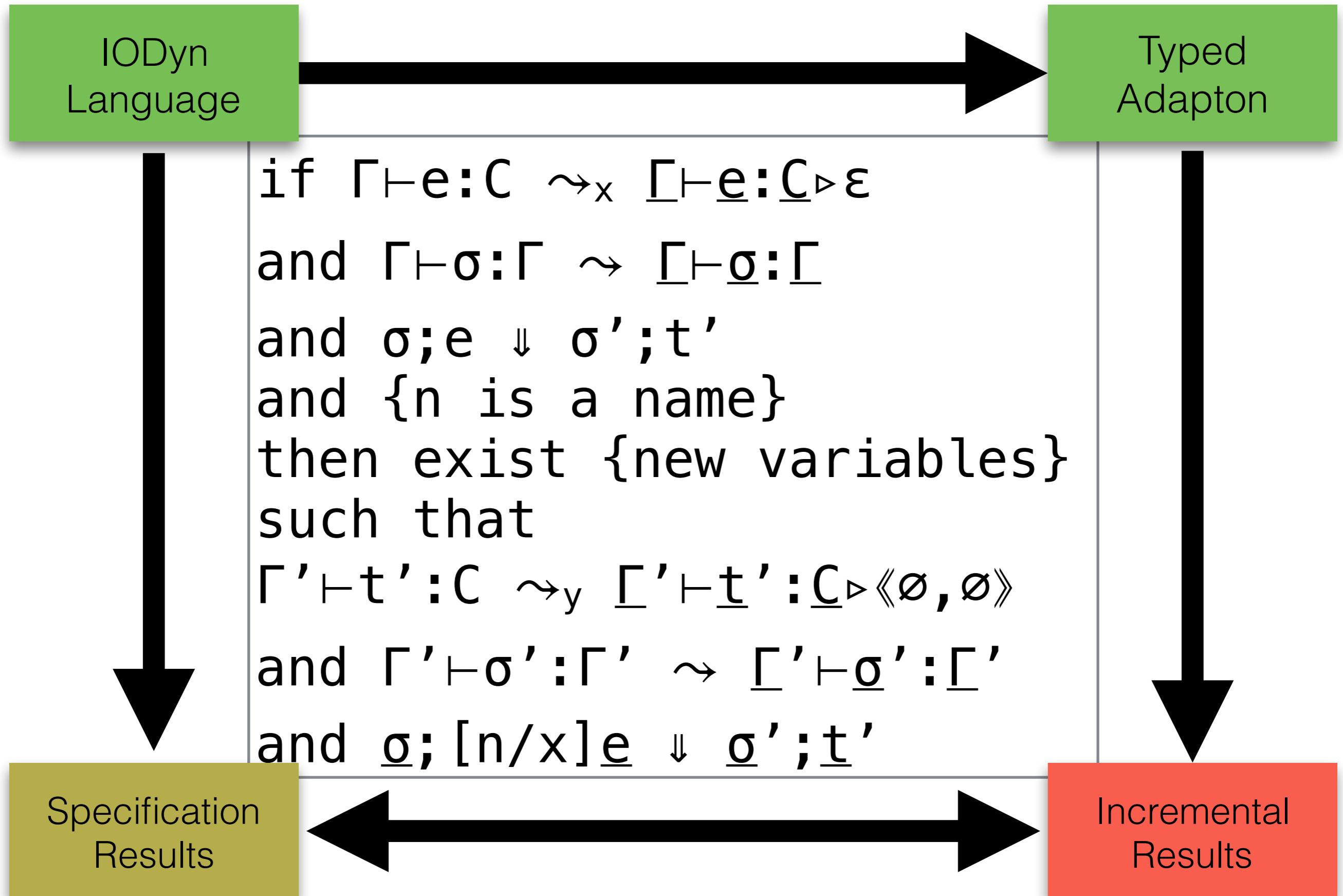
Not directly possible, instead,
use a fold, checking each token

Translation soundness

Translation soundness

if $\Gamma \vdash e : C \rightsquigarrow_x \underline{\Gamma} \vdash \underline{e} : \underline{C} \triangleright \varepsilon$
and $\Gamma \vdash \sigma : \Gamma \rightsquigarrow \underline{\Gamma} \vdash \underline{\sigma} : \underline{\Gamma}$
and $\sigma ; e \Downarrow \sigma' ; t'$
and $\{n \text{ is a name}\}$
then exist $\{\text{new variables}\}$
such that
 $\Gamma' \vdash t' : C \rightsquigarrow_y \underline{\Gamma}' \vdash \underline{t}' : \underline{C} \triangleright \langle \emptyset, \emptyset \rangle$
and $\Gamma' \vdash \sigma' : \Gamma' \rightsquigarrow \underline{\Gamma}' \vdash \underline{\sigma}' : \underline{\Gamma}'$
and $\underline{\sigma} ; [n/x] \underline{e} \Downarrow \underline{\sigma}' ; \underline{t}'$

Translation soundness



Summary

IODyn is a new incremental language that abstracts away much of the complexity of incremental code

IODyn combines the optimizations and safety of multiple previous projects

I've been working on defining and type-checking the IODyn source language

Next steps include implementing the translation and evaluating the results against the source spec

www.github.com/cuplv/iodyn-lang.rust

kyleheadley.github.io